Determinantal representations of cubic curves

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Background

Elliptic curves have profound influence in mathematics. Since ancient times they turn up in the most astonishing places, joining together algebra and geometry. Recently they have become popular in **number theory** (cryptography of elliptic curves), **optimization** (semidefinite programming SDP) and also in **theoretical physics** (mirror symmetry of elliptic curves).

Background

The abundance of results is due to the following two classical facts for smooth plane cubics:

- It can be brought by a change of coordinates into the Weierstrass canonical form, or equivalently the Hesse canonical form.
- It can be equipped by a group law (induced by the Jacobian group variety).

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Background Notation

Notation

- we work over the field \mathbb{C} , sometimes we restrict to \mathbb{R} ,
- F(x, y, z) homogeneous polynomial of degree 3,
- C a smooth curve defined by $\{F(x, y, z) = 0\} \subset \mathbb{P}^2$.

$$\begin{array}{cccc} & \text{projective plane } \mathbb{P}^2: & \text{affine plane } \mathbb{C}^2: \\ \text{points:} & (x,y,z) = & \\ & (x/z,y/z,1) & \Longleftrightarrow & (X,Y) \\ \text{curves:} & y^2z = (x+z)(x^2 + \varepsilon z^2) & Y^2 = (X+1)(X^2 + \varepsilon) \\ & y^2z = (x+z)x^2 & Y^2 = (X+1)X^2 \\ & y^2z = (x+z)(x^2 - \varepsilon z^2) & Y^2 = (X+1)(X^2 - \varepsilon) \end{array}$$

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Definition: Determinantal representation

It is very useful to represent *F* as a determinant of some matrix: Find a 3×3 matrix with linear terms

Definition

$$M(x, y, z) = xA + yB + zC$$

such that

$$\det M(x,y,z)=c\,F(x,y,z), \ \text{ for some } c\neq 0.$$

Matrix M is called a determinantal representation of C.

Clearly, multiplying a determinantal representation by invertible matrices preserves the underlying curve. Two determinantal representations *M* and *M'* equivalent if there exist $X, Y \in GL(3, \mathbb{C})$ such that

$$M' = XMY.$$

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Existence of a symmetric determinantal representation

Every cubic curve has a determinantal representation. The following theorem constructs a symmetric one:

Theorem (J. Harris, 1979, p. 696)

There exist precisely three points $(a, b) \in \mathbb{C}^2$ such that

 $aF = \operatorname{Hes}(bF + \operatorname{Hes}(F)),$

where Hes is the Hessian i.e., the determinant of the second partial derivatives matrix. The resulting three symmetric determinantal representations of F are inequivalent.

Using elementary transformations [Vinnikov] we can obtain all determinantal representations of *F* from a given one.

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Canonical form Inflection point Algorithm

Weierstrass canonical form

Theorem

By a projective change of coordinates, every irreducible curve can be brought into the Weierstrass form

$$y^2z = x^3 + pxz^2 + qz^3, \quad p,q \in \mathbb{C}$$

or equivalently $y^2 z = x(x + \theta_1 z)(x + \theta_2 z), \quad \theta_1, \theta_2 \in \mathbb{C}.$

Moreover, every reduced curve is projectively equivalent to one of the

$$x^3, x^2y, xy(x+y), xyz$$
 or
 $(\alpha x + \beta y + \gamma z)(x^2 - yz)$ for some $\alpha, \beta, \gamma \in \mathbb{C}$.

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Canonical form Inflection point Algorithm

Why do we want the Weierstrass canonical form?

Corollary

Any coordinate independent statement that holds for a Weierstrass cubic, holds for any irreducible cubic curve.

We will use this to show:

• Determinantal representations of any cubic curve *C* are in one to one correspondence with affine points on *C*.

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Canonical form Inflection point Algorithm

Inflection point

Every irreducible cubic has inflection points: $\{F = 0\} \cap \{\text{Hes } F = 0\} \subset \mathbb{P}^2.$

Proposition

If we find an inflection point on C, we can put it into the Weierstrass form.

Change the coordinates so that the inflection point is (0, 1, 0)and the inflection tangent is z = 0. Considering all possible monomials occurring in F yields the Weierstrass form.

Corollary

When the defining polynomial F is real, a real change of coordinates gives the Weierstrass form with $p, q \in \mathbb{R}$.

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Canonical form Inflection point Algorithm



- The enumerative problem of locating flexes of a plane cubic is solvable, since the corresponding Galois group is solvable [Harris, 1979].
- When C contains a rational point [Silverman and Tate, 1992] provided an algorithm that puts it into a Weierstrass form.

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Explicit Construction Definite representation Semidefinite programming

Main theorem

Theorem

Determinantal representations of $y^2 z = x(x + \theta_1 z)(x + \theta_2 z)$ are

$$M_{l,t}(x, y, z) = x \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix} + y \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} + z \begin{pmatrix} \frac{3}{4}t^2 + \frac{1}{2}t(\theta_1 + \theta_2) - \frac{1}{4}(\theta_1 - \theta_2)^2 & l & \frac{1}{2}(\theta_1 + \theta_2 + t) \\ -l & -l & -t & 0 \\ \frac{1}{2}(\theta_1 + \theta_2 + t) & 0 & -1 \end{pmatrix}$$

where $l, t \in \mathbb{C}$ and $l^2 = t(t + \theta_1)(t + \theta_2)$.

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Explicit Construction Definite representation Semidefinite programming

Real cubic

It is natural to consider the subset of cubics defined by real polynomials. Representation $M_{l,t}(x, y, z)$ in the Main theorem is self-adjoint iff *t* is real and I = is is purely imaginary. The set of all nonequivalent self-adjoint determinantal representations of C can thus be parametrised by affine points

on F(t, I) whose first coordinate is real and the second coordinate is purely imaginary.

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Explicit Construction Definite representation Semidefinite programming

Real cubic

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Explicit Construction Definite representation Semidefinite programming

Smooth cubics $-s^2 = t(t + \theta_1)(t + \theta_2)$



Explicit Construction Definite representation Semidefinite programming

Singular cubics $-s^2 = t^3$, $-s^2 = t^2(t-1)$, $-s^2 = t(t-1)^2$.



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Explicit Construction Definite representation Semidefinite programming

Definite representation

Definition

A self-adjoint determinantal representation is definite if there exists a point $(a, b, c) \in \mathbb{R}^3$ such that M(a, b, c) is positive definite.

Note:

Definitness is preserved under a real change of coordinates.

Theorem

Let C be a real cubic

$$y^2 z = x(x + \theta_1 z)(x + \theta_2 z).$$

Determinantal representations $M_{l,t}$ from the Main theorem are definite for t > 0.

Explicit Construction Definite representation Semidefinite programming

Semidefinite programming SDP

Definition

The set of points where a given self-adjoint determinantal representation representation is positive definite, $S = \{(x, y, z) \in \mathbb{R}^3 : M(x, y, z) \succeq 0\}$ is called spectrahedron and $M(x, y, z) \succeq 0$ is an LMI (linear matrix inequality) representation of S.

Theorem (Vinnikov, 2007)

Rigidly convex sets are exactly spectrahedra bounded by the determinant of some LMI representation. Spectrahedra are precisely the sets on which semidefinite programming can be performed.

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