Lower General Position Sets in Graphs

Gabriele Di Stefano [a]

gabriele.distefano@univaq.it

Sandi Klavžar [b,c,d]

sandi.klavzar@fmf.uni-lj.si

Aditi Krishnakumar [e]

aditikrishnakumar@gmail.com

James Tuite [e]

james.t.tuite@open.ac.uk

Ismael G. Yero [f]

ismael.gonzalez@uca.es

June 19, 2023

Abstract

A subset S of vertices of a graph G is a *general position set* if no shortest path in G contains three or more vertices of S. In this paper, we generalise a problem of M. Gardner to graph theory by introducing the *lower general position number* $\mathrm{gp}^-(G)$ of G, which is the number of vertices in a smallest maximal general position set of G. We show that $\mathrm{gp}^-(G) = 2$ if and only if G contains a universal line and determine this number for several classes of graphs, including Kneser graphs $K(n, 2)$, line graphs of complete graphs, and Cartesian and direct products of two complete graphs. We also prove several realisation results involving the lower general position number, the general position number and the geodetic number, and compare it with the lower version of the monophonic position number. We provide a sharp upper bound on the size of graphs with given lower general position number. Finally we demonstrate that the decision version of the lower general position problem is NP-complete.

Keywords: general position number; geodetic number; universal line; computational complexity; Kneser graphs; line graphs

AMS Subj. Class. (2020): 05C12, 05C69, 68Q25

1 Introduction

The *general position problem* originated in a puzzle by Dudeney in his book [16]. It can be stated as follows: what is the largest number of pawns that can be placed on an $n \times n$ chessboard such that no three pawns lie on a straight line? This geometrical problem is also known as the *no-three-in-line problem*. An obvious upper bound is $2n$, which is achieved for $n \leq 46$. However, for larger n the problem remains open.

Erdős showed how to place $n - o(n)$ pawns on the chessboard with no three in line (see [35]), which was subsequently improved by Hall et al. to $\frac{3n}{2} - o(n)$ pawns [21]. It is conjectured that the true answer is $\frac{\pi n}{\sqrt{3}} - o(n)$ in [20] (see [33] for a correction). The problem has been called 'one of the oldest and most extensively studied geometric questions concerning lattice points' [7].

The general position problem was generalised to the setting of graph theory independently in [9, 28, 29] as follows.

Definition 1.1. A set $S \subseteq V(G)$ is in *general position* if no shortest path in G contains three or more vertices of S; such a set is a *general position set*. The *general position number* $\mathrm{gp}(G)$ of G is the number of vertices in a largest general position set. The *general position problem* asks for a largest general position set in a given graph.

Papers in this very active field of research include [19, 27, 31, 32, 38, 40, 42]. Several variations of the problem have been considered, including using the Steiner distance instead of the regular graph distance [24], or confining attention to shortest paths of bounded length [26]. Games involving general position sets have also been treated in [10] and [25], a dynamic variant of the problem was considered in [23] and a local version of general position sets was studied in [37].

The edge version of the general position problem has also been recently studied in [30]. A related problem is the *monophonic position problem* obtained by replacing 'shortest path' in the general position problem by 'induced path', see [39]. Another variant of the general position problem is the *mutual visibility problem* that asks for a largest set of vertices S, such that for each pair of vertices in S there is a shortest path connecting them that does not contain a third vertex of S, see [15].

A new slant on this old problem was given by Martin Gardner (the modern day Dudeney), who asked the following question in his column in Scientific American: 'Instead of asking for the maximum number of counters that can be put on an order-n board, no three in line, let us ask for the minimum that can be placed such that adding one more counter on any vacant cell will produce three in line' [17]. If a greedy algorithm is used to produce a general position set, then the answer to Gardner's problem represents the worst-case output. This problem was treated in [3, 4, 14]. The most recent of these articles, [4], refers to this problem as the *geometric domination problem* and gives a lower bound of $\Omega(n^{\frac{2}{3}})$ and an upper bound of $2\left\lceil \frac{n}{2} \right\rceil$ for an $n \times n$ grid.

In this paper we extend Gardner's problem to graph theory by asking for the smallest maximal general position sets, i.e. the smallest general position sets that cannot be extended without creating three in a line.

Definition 1.2. A general position set S in a graph G is *maximal* if there is no general position set of G containing S as a proper subset. The *lower general position number* $\mathrm{gp}^-(G)$ of G is the number of vertices in a smallest maximal general position set of G, also called a *lower general position set*.

An example of these concepts in the context of the Petersen graph P can be seen in Fig. 1. From [29] we know that $\mathrm{gp}(P) = 6$; the set of white vertices in the

figure represents a largest general position set, whilst the grey vertices form a smallest maximal general position set, so that $\mathrm{gp}^-(P) = 4$.

Since any vertex of a (non-trivial) graph G is not a maximal general position set, for any graph we have:
$$2 \leq \mathrm{gp}^-(G) \leq \mathrm{gp}(G). \tag{1}$$
It is not difficult to see that for any tree T (or, more generally, any graph with a bridge) we have $\mathrm{gp}^-(T) = 2$, so it is possible to have equality with the lower bound in Inequality 1. In Section 2 we will take a closer look at the graphs G which satisfy $\mathrm{gp}^-(G) = 2$. The upper bound in Inequality 1 is also tight, since any complete graph K_n satisfies $\mathrm{gp}^-(K_n) = n = \mathrm{gp}(K_n)$.
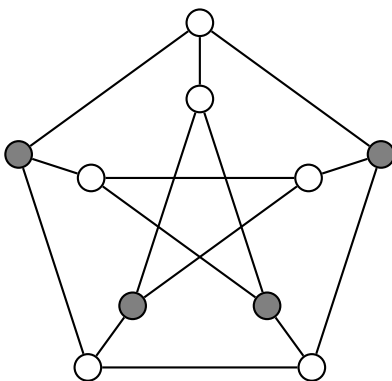

Figure 1: The Petersen graph with a maximum general position set (white) and a lower general position set (grey)

We now provide some necessary definitions. All graphs considered in this paper are finite, undirected and simple. The order and the size of a graph G will be denoted by $n(G)$ and $m(G)$, respectively. We will write $u \sim v$ if vertices u and v are adjacent. The subgraph of G induced by a subset $S \subseteq V(G)$ will be denoted by $G[S]$. The *distance* $d_G(u,v)$ between vertices u and v of a graph G is the length of a shortest u,v-path. The *interval* $I[u,v]$ of u and v is the set of vertices that lie on at least one shortest u,v-path. The *geodetic closure* $I[S]$ of a set $S \subseteq V(G)$ is the union $\bigcup_{u,v \in S} I[u,v]$ and S is *geodetic* if $I[S] = V(G)$. The *geodetic number* $\mathrm{g}(G)$ of G is the number of vertices in a smallest geodetic set of G, see [8].

The *clique number* $\omega(G)$ is the number of vertices in a largest clique (i.e. set of mutually adjacent vertices) of G, whilst $\omega^-(G)$ will represent the number of vertices in a smallest maximal clique of G. A subset $S \subseteq V(G)$ is an *independent union of cliques* if each component of $G[S]$ is a clique; the number of vertices in a largest independent union of cliques is the *independent union of cliques number* $\alpha^\omega(G)$, whereas we will denote the number of vertices in a smallest maximal independent union of at least two cliques by $\alpha^{\omega^-}(G)$ (for convenience we do not count a clique as an independent union of cliques here). The join of graphs G and H will be denoted by $G \vee H$. We also adopt the convention that $[n]$ stands for the set $\{1, \ldots, n\}$.

The plan of this paper is as follows. In Section 2 we consider graphs having lower general position number equal to two. In particular, we show that these graphs coincide with the graphs that contain a universal line. In Section 3 we prove several

realisation results involving the lower general position number, the general position number and the geodetic number. We also give an upper bound on the size of a graph with given lower general position number and characterise the graphs which attain equality. In Section 4 we prove that the decision version of the lower general position problem is NP-complete. In Section 5 we determine the lower general position number of the Kneser graphs $K(n, 2)$, the line graphs of complete graphs, and the Cartesian and direct products of two complete graphs. In Section 6 we relate the lower general position number to the lower monophonic position number. Finally in Section 7 we suggest several avenues for further research.

2   Graphs G with $\mathrm{gp}^-(G) = 2$

Based on the trivial lower bound in Inequality (1), it is natural (as with many other graph parameters) to consider characterising the family of graphs achieving equality. However, as we now show, this appears to be a very challenging problem. To see this, let us first invoke the following concept from the theory of metric spaces. Let $M = (X, d_M)$ be an arbitrary metric space and $x, y \in X$. Then the *line* $\mathcal{L}_M(x, y)$ induced by x and y is the following set of points from M:

$$\{z \in X : \ d_M(x, y) = d_M(x, z) + d_M(z, y) \text{ or } d_M(x, y) = |d_M(x, z) - d_M(z, y)|\}.$$

The line $\mathcal{L}_M(x, y)$ is *universal* if it contains the whole set X. Considering a graph G as a metric space, these definitions transfer directly to G.

Let $\ell(M)$ denote the number of distinct lines in M. Then Chen and Chvátal [12] conjectured that if $\ell(M) < |X|$, then M has a universal line. The problem remains open; a summary of what is known about it up to 2018 can be found in [13]. For the case of graphs, the Chen-Chvátal Conjecture has been verified in particular for a certain class of graphs containing both chordal graphs and distance-hereditary graphs [2], for $(q, q - 4)$-graphs [36], and for graphs with no induced house or induced cycle of length at least 5 [1].

Closely related to the Chen-Chvátal Conjecture is an open problem [34, Problem 1.2] to determine necessary and sufficient conditions for a graph to have a universal line. It turns out that the existence of a universal line in G is equivalent to $\mathrm{gp}^-(G) = 2$.

Proposition 2.1. *Let G be a graph. Then $\mathrm{gp}^-(G) = 2$ if and only if G has a universal line.*

Proof. Suppose first that $\mathrm{gp}^-(G) = 2$ and let $\{u, v\}$ be a maximal general position set of G. Then for every $w \in V(G) \setminus \{u, v\}$, the triple of vertices u, v, w lies on a shortest path. That is, one of the following equations applies: $d_G(u, v) = d_G(u, w) + d_G(w, v)$, $d_G(u, w) = d_G(u, v) + d_G(v, w)$, $d_G(v, w) = d_G(v, u) + d_G(u, w)$. This in turn implies that $\mathcal{L}_G(u, v) = V(G)$, hence G has a universal line. Conversely, if G has a universal line $\mathcal{L}_G(u, v)$, then by the same argument, $\{u, v\}$ is a maximal general position set of G. ☐

Therefore the problem of the existence of universal lines considered in [34] is equivalent to efficiently characterising connected graphs with lower general position number 2. Simple sufficient conditions for $\text{gp}^-(G) = 2$ to hold are: (i) having a bridge; (ii) $g(G) = 2$; (iii) being bipartite, see [6]. If G is a block graph, then $\text{gp}^-(G) = 2$ if and only if K_2 is a block of G, see [34, Corollary 3.4]. In the same paper, Rodríguez-Velázquez characterised Cartesian product graphs having a universal line. (For the definition of the Cartesian product and its properties we refer to the book [22].) In view of Proposition 2.1 his result can be reformulated as follows.

Theorem 2.2. [34, Theorem 4.1] *Let G and H be non-trivial, connected graphs. Then $\text{gp}^-(G \,\square\, H) = 2$ if and only if one of the following conditions holds.*

(i) G or H has a maximal general position set consisting of two adjacent vertices.

(ii) $g(G) = 2$ and $g(H) = 2$.

Other graphs with lower general position number 2 are: (i) graphs formed from C_4 and any graph H by fixing a vertex $u \in V(C_4)$ and joining it to at least one vertex in every component of H and (ii) any complete multipartite graph with a part of cardinality 2, cf. Proposition 2.3(ii).

To conclude this section we determine the cycles and complete multipartite graphs that have lower general position number equal to 2.

Proposition 2.3. *(i) If $n \geq 3$, then*

$$\text{gp}^-(C_n) = \begin{cases} 2; & n \text{ even}, \\ 3; & n \text{ odd}. \end{cases}$$

(ii) If $t \geq 2$ and $r_1 \geq r_2 \geq \cdots \geq r_t \geq 2$, then

$$\text{gp}^-(K_{r_1,\ldots,r_t}) = \min\{t, r_t\}.$$

Proof. (i) If n is even, then C_n is bipartite and hence $\text{gp}^-(C_n) = 2$. Now let n be odd. Identify the vertex set of C_n with $\mathbb{Z}_n$ in the natural manner. As $\text{gp}(C_n) = 3$ for $n = 3$ and $n \geq 5$, we have $2 \leq \text{gp}^-(G) \leq 3$ and the result will follow if we show that any set of two vertices of C_n can be extended to form a general position set of three vertices. Without loss of generality, let this set be $S = \{0, i\}$, where $i < n - i$. Then it is easily checked that if i is odd, then the set $\{0, i, \frac{n+i}{2}\}$ is in general position and if i is even, then $\{0, i, \frac{n+i-1}{2}\}$ is in general position.

(ii) Each of the partite sets of the complete multipartite graph is a maximal general position set, so it follows that $\text{gp}^-(K_{r_1,\ldots,r_t}) \leq r_t$. Suppose that S is a maximal general position set containing vertices from different partite sets. Then S can contain at most one vertex from each of the t parts; hence S induces a clique and by maximality S must contain t vertices, one from each part. Hence $\text{gp}^-(K_{r_1,\ldots,r_t}) = \min\{t, r_t\}$. $\square$

3 Realisation results

In this section we consider the relationship between the lower general position number, the geodetic number and the general position number. We begin by determining the lower general position number of the join of two graphs. Recall our convention that when finding the smallest number of vertices in a maximal independent union of cliques $\alpha^{\omega^-}(G)$ we do not allow unions consisting of one clique.

Lemma 3.1. *If G and H are graphs, then*

$$\mathrm{gp}^-(G \vee H) = \min\{\omega^-(G) + \omega^-(H), \alpha^{\omega^-}(G), \alpha^{\omega^-}(H)\}.$$

Proof. Let S be a smallest maximal general position set of $G \vee H$. Suppose that S contains vertices of both G and H; then both $S \cap V(G)$ and $S \cap V(H)$ must induce cliques, since if $u, u' \in V(G), v \in V(H)$ and $u \not\sim u'$, then u, v, u' would be a path of length two contained in S. Conversely, any such set is a clique in $G \vee H$ that cannot be extended to a larger general position set of $G \vee H$. Hence $|S| \leq \omega^-(G) + \omega^-(H)$.

Suppose now that $S \cap V(H) = \emptyset$ (the case $S \cap V(G) = \emptyset$ is symmetrical). As all vertices of $V(G)$ are at distance at most two in $G \vee H$, S is in general position in $G \vee H$ if and only if it is an independent union of cliques. However, S will only be a maximal general position set if it is a maximal independent union of at least two cliques (otherwise we could extend S by adding any vertex from H). Thus $\mathrm{gp}^-(G \vee H) \leq \alpha^{\omega^-}(G)$.

The result now follows upon taking the minimum amongst these forms of maximal general position sets in $G \vee H$. $\square$

We also use the following known lemma in our comparison of the geodetic number and lower general position number. Recall that a vertex of a graph is *simplicial* if its neighbours induce a complete subgraph.

Lemma 3.2. [11, Theorem A] *Every geodetic set of G contains all simplicial vertices of G.*

Theorem 3.3. *Let $a \geq 2$ and $b \geq 2$ be integers. Then there is a graph G with $\mathrm{gp}^-(G) = a$ and $\mathrm{g}(G) = b$ if and only if $2 \leq a \leq b$ or $4 \leq b \leq a$.*

Proof. If $a \leq b$, then by Lemma 3.1 the graph $bK_1 \vee K_{a-1}$ has $\mathrm{gp}^-(bK_1 \vee K_{a-1}) = a$. Moreover, by Lemma 3.2, any geodetic set of $bK_1 \vee K_{a-1}$ contains the b vertices of bK_1; since this set is geodetic, we obtain $\mathrm{g}(bK_1 \vee K_{a-1}) = b$. We can thus assume in the remainder of the proof that $b < a$.

If $b = \mathrm{g}(G) = 2$, then G has a universal line and hence $\mathrm{gp}^-(G) = 2$ by Proposition 2.1. Suppose that $\mathrm{g}(G) = 3$ and let $S = \{x, y, z\}$ be a geodetic set with $|S| = 3$, so that $I[x, y] \cup I[x, z] \cup I[y, z] = V(G)$. Suppose that S is not in general position, say y lies on a shortest x, z-path P in G. Then we have $I[x, y] \cup I[y, z] \subseteq I[x, z]$, so that $\{x, z\}$ would be a geodetic set, a contradiction, implying that G would have a maximal general position set of order at most three. Thus we cannot have $b < a$ if $b \in \{2, 3\}$.

Now we deal with the remaining case $4 \leq b < a$. Let X_1, X_2, Y be cliques of orders $|X_1| = |X_2| = a - 2 > 2$ and $|Y| = b - 3 \geq 1$. Let $x_1 \in X_1$ and $x_2 \in X_2$ be fixed. We form a graph $H(a, b)$ from these cliques as follows. Add all possible edges between X_1 and Y, join x_1 to every vertex of X_2 and x_2 to every vertex of X_1. Add a new vertex w and join it to every vertex of X_2. Finally we construct $G(a, b)$ by adding a vertex z to $H(a, b)$ and joining it to every other vertex, i.e. $G(a, b) = H(a, b) \vee K_1$. An example is given in Fig. 2.
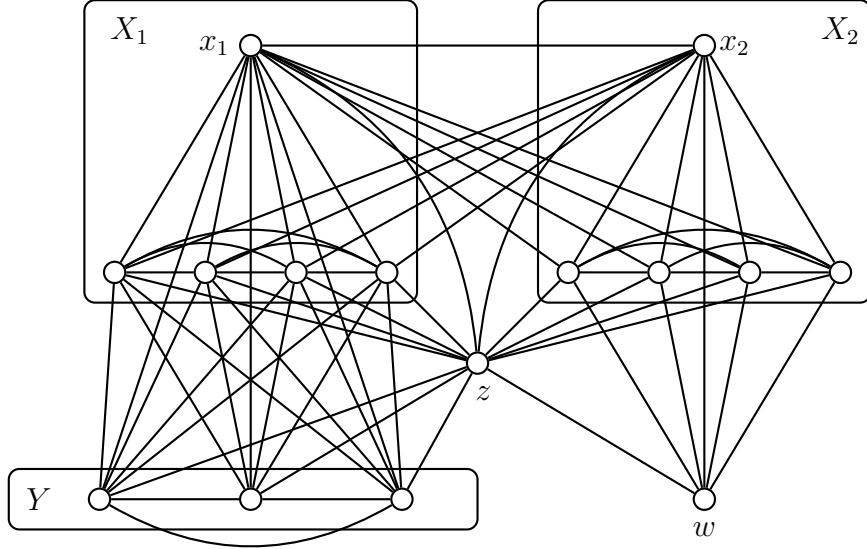

Figure 2: The graph $G(7, 6)$.

Let S be a smallest geodetic set of $G(a, b)$. Now each vertex in $Y \cup \{w\}$ is a simplicial vertex, so Lemma 3.2 yields $Y \cup \{w\} \subseteq S$, and hence $g(G(a, b)) \geq b - 2$. Furthermore $Y \cup \{x_1, x_2, w\}$ is easily seen to be a geodetic set, so that $g(G(a, b)) \leq b$, whilst if we add a single vertex to $Y \cup \{w\}$ there is no way of including both x_1 and x_2 in $I[S]$. It follows that $g(G(a, b)) = b$.

For the lower general position number, Lemma 3.1 implies that $gp^-(G(a, b)) = \min\{\omega^-(H(a, b)) + 1, \alpha^{\omega^-}(H(a, b))\}$. The set $X_2 \cup \{w\}$ induces a smallest maximal clique of $H(a, b)$, hence $\omega^-(H(a, b)) + 1 = a$. To calculate $\alpha^{\omega^-}(H(a, b))$, let S be a smallest maximal independent union of at least two cliques. Suppose that $x_1 \in S$. Observe that in this case S cannot contain vertices from both $(X_1 \setminus \{x_1\}) \cup Y \cup \{w\}$ and X_2. Hence if S contains a vertex of X_2, then S would be a clique, a contradiction, so it follows that $S \subseteq X_1 \cup Y \cup \{w\}$; as $X_1 \cup Y \cup \{w\}$ is an independent union of cliques, we would have $S = X_1 \cup Y \cup \{w\}$ and $|S| = a + b - 4 \geq a$. Thus we can assume that $x_1 \notin S$ and similarly $x_2 \notin S$.

Thus $S \subseteq (X_1 \cup X_2 \cup Y \cup \{w\}) \setminus \{x_1, x_2\}$, from which it follows that $S = (X_1 \cup X_2 \cup Y \cup \{w\}) \setminus \{x_1, x_2\}$ and $|S| = 2a + b - 8 > a$. Hence $\alpha^{\omega^-}(H(a, b)) \geq a$ and $gp^-(G(a, b)) = a$. $\qquad \square$

Note that, by Proposition 2.3(ii), for any $2 \leq a \leq b$ the complete multipartite graph with one part of cardinality b and $a - 1$ parts of cardinality a has $gp^-(G) = a$, $gp(G) = b$ and order $b + (a - 1)a$. Also the graph $bK_1 \vee K_{a-1}$ from the proof of

Theorem 3.3 has $\text{gp}^-(bK_1 \vee K_{a-1}) = a$ and $\text{gp}(bK_1 \vee K_{a-1}) = b$, but a smaller order $b + a - 1$. For $a = b$, trivially K_a is the smallest possible such graph. We next show that for $2 \leq a < b$ a smaller graph that fulfills the same conditions can be constructed.

Theorem 3.4. *For any $2 \leq a < b$ there is a graph G with $\text{gp}^-(G) = a$, $\text{gp}(G) = b$, and $n(G) = b + \max\{1, a - \lfloor \frac{b}{2} \rfloor\}$.*

Proof. Let G be the join of $K_c \cup K_{b-a+c}$ with K_{a-c}. Choose c such that $a \leq b - a + 2c \leq b$; then $\text{gp}^-(G) = a$ and $\text{gp}(G) = b$. The smallest possible choice of c is $\max\{1, a - \lfloor \frac{b}{2} \rfloor\}$. $\qquad\square$

Finding the smallest graph with $\text{gp}^-(G) = a$ and $\text{gp}(G) = b$ remains an open problem. To conclude this section we prove the following bound on the size of a graph with given lower general position number.

Theorem 3.5. *If G is a graph with $\text{gp}^-(G) = k \geq 2$ and $n(G) \geq 2k - 1$, then*

$$m(G) \leq \binom{n}{2} - k + 1.$$

Moreover, the unique graph that meets this bound is formed from a clique by deleting $k - 1$ edges adjacent to a fixed vertex.

Proof. Let $n \geq 2k - 1$ and take a clique K_n. Choose a set S of k vertices of K_n and for some vertex $u \in S$ delete all edges between u and $S \setminus \{u\}$ to form the graph $G(k)$. We claim that $\text{gp}^-(G(k)) = k$. Let S' be an arbitrary maximal general position set of $G(k)$. If $u \notin S'$, then S' must be a clique with order $n - 1 \geq k$, so suppose that $u \in S'$. If $(S \setminus \{u\}) \cap S' = \emptyset$, then S' is the clique induced by $G(k) \setminus (S \setminus \{u\})$, which contains $n - k + 1 \geq k$ vertices. If S' contains a vertex v of $S \setminus \{u\}$, then $S' \cap V(G(k) \setminus S) = \emptyset$, as any vertex $w \in G(k) \setminus S$ is contained in a shortest path u, w, v. The set S is clearly in general position, so it follows that $\text{gp}^-(G(k)) = k$.

In fact we can show that this is the only graph with size $\binom{n}{2} - k + 1$ that has lower general position number k. Assume that G is any such graph and let M be the set of $k - 1$ edges deleted from K_n to obtain G, i.e. $G = K_n \setminus M$. Let H be the subgraph of K_n induced by M. Also, let K be the set of vertices incident to the edges of M. Observe that $|K| \leq 2k - 2$. If H is connected, then $|K| \leq k$, and if H is not connected, then there is a component of H with order at most $k - 1$; in either case, let H' be a component of H with minimum order. Let u and v be two vertices of H', so that u and v are non-adjacent in G.

Let S be a maximal general position set of G containing $\{u, v\}$. S cannot contain any vertex outside H', so H must be connected, for otherwise we have exhibited a maximal general position set of G of order at most $k - 1$. As $\text{gp}^-(G) = k$ and $|K| \leq k$, we must have that $|K| = k$, $S = K$ and H is a tree. If the diameter of H is at least three, then there would be a path u_0, u_1, u_2, u_3 in H, so that u_0, u_3, u_1 would be a shortest path in G contained in S, which is impossible. We conclude that H must be a star.

8

Suppose now that there is a graph G with $m(G) \geq \binom{n}{2} - k + 2$ and $\text{gp}^-(G) = k$. Hence G can be formed by deleting a set A of at most $k - 2$ edges from K_n. Let $u \sim v$ be any edge of A. Let A_1 be the set of edges of A incident with u (apart from $u \sim v$) and A_2 be the set of edges of A incident with v (again apart from $u \sim v$). Let B_1 be the set of endvertices of the edges in A_1 apart from u, and let B_2 be the set of endvertices of the edges of A_2 apart from v. There can be overlap between the sets B_1 and B_2, but as we have deleted at most $k - 2$ edges we have $|B_1 \cup B_2 \cup \{u, v\}| \leq k - 1$. Let S be a maximal general position set of G containing the set $\{u, v\}$. S cannot contain any vertex of $V(G) \setminus (B_1 \cup B_2 \cup \{u, v\})$ and so $S \subseteq B_1 \cup B_2 \cup \{u, v\}$; this shows that G contains a maximal general position set containing at most $k - 1$ vertices, a contradiction. □

4 Computational complexity

In Section 2 we have seen that $\text{gp}^-(G) = 2$ if and only if G has a universal line. The problem of characterising graphs that have universal lines is difficult and was extensively investigated in [34]. In this section we complement these investigations by proving that the lower general position problem is NP-complete. To this end, we formally define the decision version of the problem:

Definition 4.1. LOWER GENERAL POSITION
INSTANCE: A graph G, a positive integer $k \leq n(G)$.
QUESTION: Is there a lower general position set S for G such that $|S| \leq k$?

The problem is hard to solve, as shown by the following theorem.

Theorem 4.2. *The* LOWER GENERAL POSITION *problem is NP-complete.*

Proof. Let us observe that LOWER GENERAL POSITION is in NP since, given a set of vertices S, it can be tested in polynomial time if (1) S is in general position, (2) if it is maximal and (3) if its cardinality is less than a given integer.

We prove that INDEPENDENT DOMINATING SET polynomially reduces to LOWER GENERAL POSITION. Recalling that a dominating set of a graph G is a subset V' of $V(G)$ such that for all $u \in V(G) \setminus V'$ there is a $v \in V'$ such that $uv \in E(G)$, we provide a formal definition of the decision version of the problem:

> An instance of INDEPENDENT DOMINATING SET is given by a graph G and a positive integer $k \leq n(G)$. The INDEPENDENT DOMINATING SET problem asks whether G contains a dominating set $K \subseteq V(G)$, of cardinality k or less, which is also independent.

The NP-completeness of INDEPENDENT DOMINATING SET is reported in [18]. We polynomially transform an instance (G, k) of INDEPENDENT DOMINATING SET to an instance (G', k') of LOWER GENERAL POSITION. In particular, given (G, k) we must construct a graph G' and a positive integer k' such that G has an independent

dominating set of cardinality k or less if and only if G' has a lower general position set of cardinality k' or less.

Given an instance (G, k), the graph G' is built as follows:

$$G' = (\overline{G} \cup K_{n(G)+1}) \vee K_1$$

The graph G' is thus the join of graphs $G'' = \overline{G} \cup K_{n(G)+1}$ and K_1, where G'' is the disjoint union of the complement of G and a clique graph with $n(G)+1$ vertices. Then the construction of G' can be achieved in polynomial time. As for k', we set $k' = k+1$.

We have to show that an instance (G, k) of INDEPENDENT DOMINATING SET has a positive answer if and only if (G', k'), the corresponding instance of LOWER GENERAL POSITION, has a positive answer.

By Lemma 3.1, the lower general position number of G' is given by

$$\mathrm{gp}^-(G') = \min\{\omega^-(\overline{G} \cup K_{n(G)+1}) + \omega^-(K_1), \alpha^{\omega^-}(\overline{G} \cup K_{n(G)+1}), \alpha^{\omega^-}(K_1)\}.$$

Considering that $\omega^-(K_1) = 1$, that $\alpha^{\omega^-}(\overline{G} \cup K_{n(G)+1}) \geq \omega^-(\overline{G} \cup K_{n(G)+1}) + 1$ and that $\alpha^{\omega^-}(K_1)$ is not defined, we have that

$$\mathrm{gp}^-(G') = \omega^-(\overline{G} \cup K_{n(G)+1}) + 1 = \omega^-(\overline{G}) + 1. \tag{2}$$

Moreover, observe that a smallest maximal clique in a given graph is an independent dominating set in the complement graph. Indeed, a clique is an independent set in the complement graph and, as the clique is maximal, any vertex not in the independent set of the complement graph must be adjacent to a vertex of the independent set, so the independent set is also dominating.

Assume that the instance (G, k) of INDEPENDENT DOMINATING SET has a positive answer. Then in $\overline{G}$ there is a maximal clique having cardinality at most k. Since $k \leq n(G)$, then $\omega^-(\overline{G} \cup K_{n(G)+1}) \leq k$ and hence by (2), $\mathrm{gp}^-(G') = \omega^-(\overline{G} \cup K_{n(G)+1}) + 1 \leq k + 1 = k'$, which in turn means that (G', k'), the instance of LOWER GENERAL POSITION, has a positive answer.

Assume now that there is a maximal general position set S in G' with cardinality k' or less, that is $|S| \leq k'$. By (2), if S contains a vertex in $K_{n(G)+1}$, then, as $|S|$ is maximal, it must contain all the vertices in $K_{n(G)+1}$ and at least one vertex in $\overline{G} \vee K_1$. So the cardinality of S is at least $n(G) + 2 \geq k' + 1$, a contradiction. Thus S does not contain vertices in $K_{n(G)+1}$. It now follows from Lemma 3.1 that S consists of the vertex of K_1 and a smallest maximal clique in $\overline{G}$, implying that there exists an independent dominating set in G of cardinality $|S| - 1 \leq k$. $\square$

5 Lower general position number of some families

In this section we determine the lower general position number of the Kneser graphs $K(n, 2)$, the line graphs of complete graphs, and the Cartesian and direct products of two complete graphs.

Recall that the Kneser graph $K(n, 2)$ is the graph with vertex set consisting of all subsets of cardinality two of the set $[n]$, with an edge between two such subsets if and only if they are disjoint.

Theorem 5.1. *If $n \geq 3$, then*

$$\mathrm{gp}^-(K(n, 2)) = \begin{cases} 3; & n \in \{3, 6, 7\}, \\ 4; & n \in \{5, 8, 9\}, \\ 5; & n \in \{10, 11\}, \\ 6; & n = 4 \ or \ n \geq 12. \end{cases}$$

Proof. The result is trivial for $n = 3, 4$. The case $n = 5$ is the Petersen graph. Hence we can assume that $n \geq 6$. We divide our argument into three cases depending on the structure of the general position set. By the result of [5], any general position set is either a clique or an independent union of cliques (with a couple of additional technical properties). Let K be any maximal general position set of $K(n, 2)$.

Case 1: The subgraph induced by K contains a clique W of order ≥ 3.
Suppose in this case that the subgraph induced by K contains a component W' apart from W and let $\{a, b\}$ be a vertex of W'. Without loss of generality, let three vertices of W be $\{1, 2\}$, $\{3, 4\}$ and $\{5, 6\}$. Then the set $\{a, b\}$, identifying a vertex of W', must have a non-empty intersection with each of these sets $\{1, 2\}$, $\{3, 4\}$ and $\{5, 6\}$, which is impossible. Thus K consists of a single clique W, which is maximal if and only if it contains $\lfloor \frac{n}{2} \rfloor$ sets. Thus in this case the maximal general position sets have order $\lfloor \frac{n}{2} \rfloor$.

Case 2: K is an independent set.
As $K(n, 2)$ has diameter two any independent set is in general position. By the Erdős-Ko-Rado Theorem the independence number of $K(n, 2)$ is $n-1$. Let $\{1, 2\}$ be a vertex of K. Any set containing only vertices of the form $\{1, i\}$ is independent (similarly for a set containing only vertices of the form $\{2, i\}$), but the only maximal such set has cardinality $n - 1$, as given by the Erdős-Ko-Rado Theorem, which is greater than the maximal general position sets considered in Case 1 (unless $n = 5$, in which case $K(5, 2)$ is triangle-free and an independent set of order 4 is best possible). The only other maximal independent sets of $K(n, 2)$ have the structure $\{\{1, 2\}, \{1, 3\}, \{2, 3\}\}$; However, this is not a maximal general position set, as the set $\{1, 4\}$ can be added to make a larger general position set.

Case 3: K induces a clique K_2.
Without loss of generality, assume that the clique in question is $\{\{1, 2\}, \{3, 4\}\}$. Then any other vertex of K must be a subset of $[4]$. The set of 6 subsets of cardinality 2 of $[4]$ does form a maximal general position set isomorphic to $3K_2$. In this case the order of the maximal general position set is thus 6; this is a smallest possible maximal general position set for $n \geq 12$. $\square$

We continue with the line graphs of complete graphs, denoted by $L(K_n)$.

Theorem 5.2. *If $n \geq 2$, then*

$$\text{gp}^-(L(K_n)) = \begin{cases} \frac{n}{2}; & n \text{ even,} \\ \frac{n+3}{2}; & n \text{ odd.} \end{cases}$$

Proof. The case $n = 2$ is trivial, so we assume that $n \geq 3$. Let S be a lower general position set of $L(K_n)$. It is shown in [19] that the vertices of a maximal general position set of $L(K_n)$ correspond to edges in K_n that induce a disjoint union of triangles and stars containing either one or at least three edges. Let r be the number of triangles in S. Observe that if n is even, then a perfect matching in K_n is a maximal general position set, whilst if n is odd the disjoint union of a triangle and a matching of cardinality $\frac{n-3}{2}$ is a maximal general position set. This shows that $\text{gp}^-(L(K_n)) \leq \frac{n}{2}$ if n is even and $\text{gp}^-(L(K_n)) \leq \frac{n+3}{2}$ if n is odd.

Suppose that each edge of K_n in S lies in a triangle (in particular this requires $3|n$); then $|S| = n$, which is larger than the maximal general position sets previously constructed. Thus, we can assume that S contains at least one star. Suppose that there is a vertex of K_n that is not covered by an edge of S; since this vertex could be joined to the centre of a star to form a larger general position set, this would contradict the maximality of S. It follows that $|S| = n - r$. If n is even, it follows that $r \leq \frac{n}{2}$, so that $|S| \geq \frac{n}{2}$. Now suppose that n is odd. If S contains no triangles, then since S cannot contain an induced star of order 3 in K_n, S must contain a star on at least five vertices of K_n, in which case $r \leq \frac{n-3}{2}$ and $|S| \geq \frac{n+3}{2}$. Thus we can assume that S contains a triangle, in which case we again have $r \leq \frac{n-3}{2}$. $\square$

We finally consider the Cartesian product of two complete graphs K_r and K_s with $r, s \geq 2$.

Theorem 5.3. *If $r, s \geq 2$, then*

$$\text{gp}^-(K_r \square K_s) = \min\{r, s\}.$$

Proof. Since the vertices of any copy of K_r or K_s in $K_r \square K_s$ form a maximal general position set, trivially $\text{gp}^-(K_r \square K_s) \leq \min\{r, s\}$. Now, assume that there is a maximal general position set S of $K_r \square K_s$ of cardinality smaller than $\min\{r, s\}$. First note that S cannot be a proper subset of the vertex set of any copy of K_r or of K_s. Also, there must exist a copy of K_r, say $K_r^{(i)}$, and a copy of K_s, say $^{(j)}K_s$, which do not contain any vertex of S. This immediately allows us to observe that, independently of the structure of S, the unique vertex of $K_r \square K_s$ belonging to both copies $K_r^{(i)}$ and $^{(j)}K_s$, together with the set S would also form a general position set of $K_r \square K_s$. Therefore, S is not maximal, which leads to the desired equality. $\square$

Theorem 5.3 should be compared with [19, Theorem 3.2] which asserts that if $r, s \geq 2$, then $\text{gp}(K_r \square K_s) = r + s - 2$.

Continuing the theme of products of complete graphs, the direct product is also of interest. Given two graphs G and H, the direct product of G and H is the graph $G \times H$ whose vertex set is $V(G) \times V(G)$ and two vertices $(g, h), (g', h') \in V(G \times H)$

are adjacent in $G \times H$ if and only if $gg' \in E(G)$ and $hh' \in E(H)$. For any $g \in V(G)$ we call the subgraph of $G \times H$ induced by $\{g\} \times H$ a *H-layer* of $G \times H$ (and similarly if $h \in V(H)$ the subgraph induced by $G \times \{h\}$ is a *G-layer*). Note that, unlike the case of Cartesian products, a G-layer will be isomorphic to G only if G is an empty graph.

Theorem 5.4. *If $r, s \geq 2$ with $(r, s) \neq (2, 2)$, then*

$$\mathrm{gp}^-(K_r \times K_s) = \min\{r, s, 4\}.$$

Proof. We can assume without loss of generality that $2 \leq r \leq s$ and $s \geq 3$. A subset W of $V(K_r \times K_s)$ induces a clique if and only if all vertices of W lie in distinct layers of $K_r \times K_s$, whereas a subset $A \subseteq V(K_r \times K_s)$ is independent if and only if it lies within a single layer.

We show first that any maximum clique W is a maximal general position set. Let $(x, y) \in V(K_r \times K_s) \setminus W$; we must show that $W \cup \{(x, y)\}$ is not in general position. If $r \geq 3$, then W contains vertices (x, j_1) and (i_2, j_2) such that $i_2 \neq x$ and $j_2 \neq y$; then $(x, y), (i_2, j_2), (x, i_1)$ is a shortest path and (x, y) cannot be added to W to make a larger general position set. If $r = 2$, let $W = \{(i_1, j_1), (i_2, j_2)\}$. Without loss of generality, $x = i_1$. If $y \neq j_2$, then $(i_1, j_1), (i_2, j_2), (i_1, y)$ is a shortest path, whereas if $j_2 = y$, then for $j \in V(K_s) \setminus \{j_1, j_2\}$ the path $(i_2, j_2), (i_1, j_1), (i_2, j), (i_1, j_2) = (x, y)$ is a geodesic. This shows that $\mathrm{gp}^-(K_r \times K_s) \leq \min\{r, s\}$. For $r = 2$, this yields $\mathrm{gp}^-(K_2 \times K_s) = 2 = \min\{2, s, 4\}$ and $K_2 \times K_s$ has a universal line. For $r = 3$, this yields $\mathrm{gp}^-(K_3 \times K_s) \leq 3$.

For $r \geq 3$ any layer is a maximal general position set. Suppose that $\{i\} \times K_s$ is a K_s-layer (the argument for K_r-layers is identical). The vertices of $\{i\} \times K_s$ lie at distance 2 from each other, so the set is in general position, whilst if we add any vertex (i', j), $i' \neq i$, then for $j_1, j_2 \in V(K_s) \setminus \{j\}$ the path $(i, j_1), (i', j), (i, j_2)$ is a geodesic.

Therefore, as any pair of vertices either forms an independent set or a clique, for $r \geq 3$ any set of two vertices of $K_r \times K_s$ can be extended to a larger general position set and $K_r \times K_s$ does not have a universal line, so that $\mathrm{gp}^-(K_r \times K_s) \geq 3$. In particular, it follows that $\mathrm{gp}^-(K_3 \times K_s) = 3 = \min\{3, s, 4\}$ and we can assume that $r \geq 4$.

Consider the set $S = \{(i, j), (i, j'), (i', j), (i', j')\}$ with $i \neq i'$ and $j \neq j'$. First observe that the edges $(i, j)(i', j'), (i, j')(i', j)$ are in $E(K_r \times K_s)$, and that none of these four vertices belongs to a shortest path between two of the remaining ones. Thus, S is a general position set of $K_r \times K_s$. Moreover, any other vertex (x, y) of $K_r \times K_s$ not in S belongs to a shortest path between two vertices of S (without loss of generality $x \neq i$ and $y \neq j$, so (x, y) lies on a path of length 2 between either (i, j) and (i, j') or (i, j) and (i', j)). Consequently, S is a maximal general position set of $K_r \times K_s$, which leads to $3 \leq \mathrm{gp}^-(K_r \times K_s) \leq 4$.

Suppose that $K_r \times K_s$ has a maximal general position set S of order 3. As $r \geq 4$, by the preceding argument if S is a clique or an independent set, then S could be extended to a larger set, so, since a general position set is an independent union of cliques [5], S must induce a graph isomorphic to $K_1 \cup K_2$ (a graph on three vertices with only one edge). If (i, j) and (i', j') are the vertices of the induced K_2 in S, then the vertex

13

corresponding to the K_1 must be either (i', j) or (i, j'), say (i', j); however, in this case S could be extended by adding (i, j') to give a general position set of order 4 with the form discussed above. Therefore for $r \geq 4$ we have $\mathrm{gp}^-(K_r \times K_s) = 4 = \min\{r, s, 4\}$ and the result is proven. $\qquad \square$

6 Connection with lower monophonic position number

In this section we relate the lower general position number to the monophonic position number mentioned in Section 1. The monophonic position number was introduced in [39] as follows. A path P in a graph G is *induced* or *monophonic* if G contains no chords between non-consecutive vertices of P. A set $M \subseteq V(G)$ is in *monophonic position* if no induced path in G contains three or more vertices of M; the *monophonic position number* $\mathrm{mp}(G)$ of G is the number of vertices in a largest monophonic position set. It was shown in [39] that for any graph G we have $\mathrm{mp}(G) \leq \mathrm{gp}(G)$ and that for any $2 \leq a \leq b$ there exists a graph with $\mathrm{mp}(G) = a$ and $\mathrm{gp}(G) = b$. This question was explored further in [41], which asked for the smallest possible order of a graph G with $\mathrm{mp}(G) = a$ and $\mathrm{gp}(G) = b$ for given $a \leq b$.
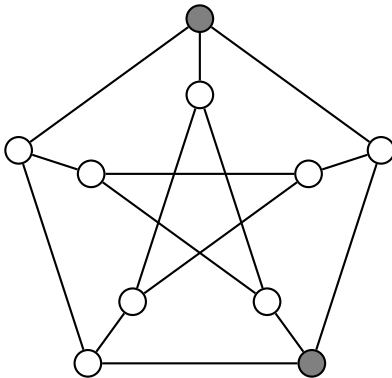

Figure 3: The Petersen graph with a lower monophonic position set (grey)

By analogy with the lower general position number, we define the *lower monophonic position number* $\mathrm{mp}^-(G)$ of G to be the number of vertices in a smallest maximal monophonic position set of G. For any graph G with order $n \geq 2$ we have $\mathrm{mp}^-(G) \geq 2$ and it is easily verified that the construction of Theorem 3.4 shows that for any $2 \leq a \leq b$ there is a graph G with $\mathrm{mp}^-(G) = a$ and $\mathrm{mp}(G) = b$. For an example of this concept, see Figure 3, which displays a lower monophonic position set in the Petersen graph P. Recall that $\mathrm{gp}^-(P) = 4$, so that in this case $\mathrm{mp}^-(P) < \mathrm{gp}^-(P)$. Intuition might suggest that the relation $\mathrm{mp}^-(G) \leq \mathrm{gp}^-(G)$ holds generally, as with the 'ordinary' general and monophonic position numbers. Interestingly, this turns out to be false, as demonstrated by the graphs in Figure 4. It remains an open question for which values of a, b there exists a graph with $\mathrm{mp}^-(G) = a$ and $\mathrm{gp}^-(G) = b$ when $a > b$. However, we were able to prove the following realisation result.

Theorem 6.1. *For any $2 \leq a \leq b$ there exists a graph G with $\mathrm{mp}^-(G) = a$ and*
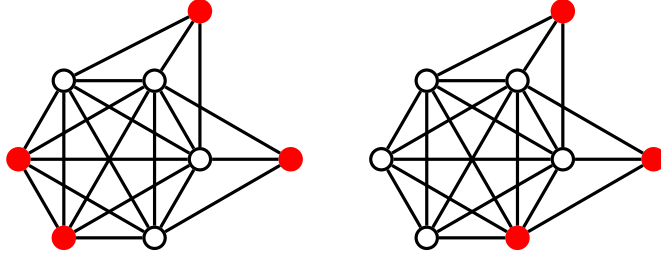
Figure 4: A graph G with $\text{mp}^-(G) = 4$ (left) and $\text{gp}^-(G) = 3$ (right)

$\text{gp}^-(G) = b$.

Proof. If $a = b$ then obviously the clique K_a will suffice, so assume that $b > a$. Take a cycle of length 6 and identify its vertex set with $\mathbb{Z}_6$ in the natural way. Now expand each vertex i into a clique W_i and add all possible edges between W_i and W_{i+1} for $0 \le i \le 5 \pmod 6$.

Suppose that a is even. Choose the cliques such that $|W_0| = |W_3| = \frac{a}{2}$ and all other cliques have order $b - \frac{a}{2}$. In Fig. 5 the construction for $a = 4$ and $b = 6$ is presented.
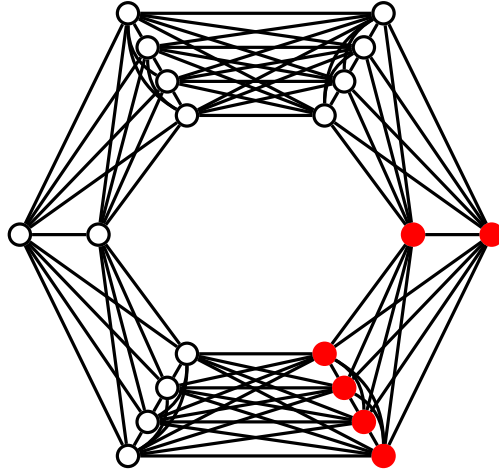

Figure 5: Construction for $a = 4$ and $b = 6$: the lower gp-set

As the monophonic position number of a cycle of length 6 is 2, it follows that the lower mp-number of this graph is the sum of the orders of the two smallest cliques; as $b \ge a$ the lower mp number is a. Similarly it follows by Proposition 2.3 that the lower gp number is the smallest element in the set

$$\{|W_i| + |W_{i+1}| : 0 \le i \le 5 \pmod 6\} \cup \{|W_0| + |W_2| + |W_4|, |W_1| + |W_3| + |W_5|\},$$

which is b. Therefore this graph has the required parameters.

If a is odd it can be verified that the graph with $|W_0| = \frac{a+1}{2}$, $|W_3| = \frac{a-1}{2}$ and all other parts of size $b - \frac{a-1}{2}$ works. $\qquad\square$

7 Concluding remarks

In this paper we considered general position sets of smallest cardinality that are maximal with respect to the set inclusion property. We conclude by mentioning some promising directions for future research suggested by our results.

- In Section 6 we exhibited a graph with $\mathrm{mp}^-(G) > \mathrm{gp}^-(G)$. Characterising the pairs a, b with $a > b$ such that there exists a graph satisfying $\mathrm{mp}^-(G) = a$ and $\mathrm{gp}^-(G) = b$ remains an open problem.

- For $2 \leq a < b$, what is the smallest possible order of a graph with $\mathrm{mp}^-(G) = a$ and $\mathrm{gp}^-(G) = b$?

- In connection with inequality (1) and the difficult problem of characterising graphs G with $\mathrm{gp}^-(G) = 2$, it could be interesting to determine further families of graphs that satisfy this property.

- Also in connection with inequality (1), is it possible to characterise all graphs G with $\mathrm{gp}^-(G) = \mathrm{gp}(G)$?

- Most of the graphs studied in Section 5 have diameter two. The general position numbers of graphs with diameter two were determined in [5]. This suggests studying lower general position sets of graphs of diameter two in general.

- Cartesian products with universal lines were characterised in [34]. It would therefore be of interest to study the value of the lower general position numbers of Cartesian products.

- As discussed in Section 1, there are several noteworthy variations of the general position number in the literature, including the mutual visibility number [15], d-position sets [26], vertex position numbers [37], Steiner position numbers [24], edge general position numbers [30], mobile position sets [23], etc. We suggest studying lower versions of these parameters.

Acknowledgements

References

[1] P. Aboulker, L. Beaudou, M. Matamala, J. Zamora, Graphs with no induced house nor induced hole have the de Bruijn-Erdős property, J. Graph Theory 100 (2022) 638–652.

[2] P. Aboulker, M. Matamala, P. Rochet, J. Zamora, A new class of graphs that satisfies the Chen-Chvátal conjecture, J. Graph Theory 87 (2018) 77–88.

[3] M.A. Adena, D.A. Holton, P.A. Kelly, Some thoughts on the no-three-in-line problem, Lecture Notes in Math. 403 (1974) 6–17.

[4] O. Aichholzer, D. Eppstein, E.M. Hainzl, Geometric dominating sets-a minimum version of the No-Three-In-Line Problem, Comput. Geom. 108 (2023) Paper 101913.

[5] B.S. Anand, U. Chandran S.V., M. Changat, S. Klavžar, E.J. Thomas, Characterization of general position sets and its applications to cographs and bipartite graphs, Appl. Math. Comput. 359 (2019) 84–89.

[6] L. Beaudou, A. Bondy, X. Chen, E. Chiniforooshan, M. Chudnovsky, V. Chvátal, N. Fraiman, Y. Zwols, A De Bruijn-Erdős theorem for chordal graphs, Electron. J. Combin. 22 (2015) Paper 1.70.

[7] P. Brass, W. Moser, J. Pach, Packing lattice points in subspaces. Research Problems in Discrete Geometry. Springer, New York (2005) 417–421.

[8] U. Chandran S.V., M.C. Dourado, M.G.S. Thankachy, Computational and structural aspects of the geodetic and the hull numbers of shadow graphs, Discrete Appl. Math. 319 (2022) 499–510.

[9] U. Chandran S.V., G.J. Parthasarathy, The geodesic irredundant sets in graphs, Int. J. Math. Combin. 4 (2016) 135–143.

[10] U. Chandran S.V., S. Klavžar, P.K. Neethu, R. Sampaio, The general position avoidance game and hardness of general position games, arXiv:2205.03526 [math.CO].

[11] G. Chartrand, P. Zhang, Extreme geodesic graphs, Czechoslovak Math. J. 52 (2002) 771–780.

[12] X. Chen, V. Chvátal, Problems related to a De Bruijn-Erdős theorem, Discrete Appl. Math. 156 (2008) 2101–2108.

[13] V. Chvátal, A de Bruijn-Erdős theorem in graphs? Graph theory–favorite conjectures and open problems. 2. Probl. Books in Math., Springer, Cham, (2018) 149–176.

[14] A.S. Cooper, O. Pikhurko, J.R. Schmitt, G.S. Warrington, Martin Gardner's minimum no-3-in-a-line problem, Amer. Math. Monthly 121 (2014) 213–221.

[15] G. Di Stefano, Mutual visibility in graphs, Appl. Math. Comput. 419 (2022) 126850.

[16] H.E. Dudeney, Amusements in Mathematics, Nelson, Edinburgh (1917).

[17] M. Gardner, Mathematical games: combinatorial problems, some old, some new and all newly attacked by computer, Sci. Amer. 235 (1976) 131–137.

[18] M.R. Garey, D.S. Johnson, Computers and Intractability: A Guide to the Theory of NP-Completeness, Freeman and Co. (1979).

[19] M. Ghorbani, H.R. Maimani, M. Momeni, F.R. Mahid, S. Klavžar, G. Rus, The general position problem on Kneser graphs and on some graph operations, Discuss. Math. Graph Theory 41 (2021) 1199–1213.

[20] R.K. Guy, P.A. Kelly, The no-three-in-line problem, Can. Math. Bull. 11 (1968) 527–531.

[21] R.R. Hall, T.H. Jackson, A. Sudbery, K. Wild, Some advances in the no-three-in-line problem, J. Comb. Theory Ser. A. 18 (1975) 336–341.

[22] W. Imrich, S. Klavžar, D.F. Rall, Topics in Graph Theory. Graphs and their Cartesian Croduct, A K Peters, Ltd., Wellesley, MA, (2008).

[23] S. Klavžar, A. Krishnakumar, J. Tuite, I. Yero, Traversing a graph in general position, Bull. Aust. Math. Soc. (2023) 1–13.

[24] S. Klavžar, D. Kuziak, I. Peterin, I.G. Yero, A Steiner general position problem in graph theory, Comput. Appl. Math. 40 (2021) 1–15.

[25] S. Klavžar, P.K. Neethu, U. Chandran S.V., The general position achievement game played on graphs, Discrete Appl. Math. 317 (2022) 109–116.

[26] S. Klavžar, D.F. Rall, I.G. Yero, General d-position sets. Ars Math. Contemp. 21 (2021) Paper P1.03.

[27] S. Klavžar, I.G. Yero, The general position problem and strong resolving graphs, Open Math. 17 (2019) 1126–1135.

[28] J. Körner, On the extremal combinatorics of the Hamming space, J. Comb. Theory Ser. A 71 (1995) 112–126.

[29] P. Manuel, S. Klavžar, A general position problem in graph theory, Bull. Aust. Math. Soc. 98 (2018) 177–187.

[30] P. Manuel, R. Prabha, S. Klavžar, The edge general position problem, Bull. Malays. Math. Sci. Soc. 45 (2022) 2997–3009.

[31] P.K. Neethu, U. Chandran S.V., M. Changat, S. Klavžar, On the general position number of complementary prisms, Fund. Inform. 178 (2021) 267–281.

[32] B. Patkós, On the general position problem on Kneser graphs, Ars Math. Contemp. 18 (2020) 273–280.

[33] E. Pegg Jr., Math Games. Chessboard Tasks. `https://www.mathpuzzle.com/MAA/36-Chessboard%20Tasks/mathgames_04_11_05.html`, accessed on February 18, 2023.

[34] J.A. Rodríguez-Velázquez, Universal lines in graphs, Quaest. Math. 45 (2022) 1485–1500.

[35] K.F. Roth, On a problem of Heilbronn, J. London Math. Soc. 26 (1951) 198–204.

[36] R. Schrader, L. Stenmans, A de Bruijn-Erdős theorem for $(q, q-4)$-graphs, Discrete Appl. Math. 279 (2020) 198–201.

[37] M. Thankachy, Maya, U. Chandran S.V., J. Tuite, E. Thomas, G. Di Stefano, G. Erskine, On the vertex position number of graphs, Discuss. Math. Graph Theory (In Press) (2023) arXiv preprint arXiv:2209.00359

[38] E.J. Thomas, U. Chandran S.V., Characterization of classes of graphs with large general position number, AKCE Int. J. Graphs Comb. (2020) 1–5.

[39] E.J. Thomas, U. Chandran S.V., J. Tuite, G. Di Stefano, On monophonic position sets of graphs, Discrete Appl. Math. (2023) doi.org/10.1016/j.dam.2023.02.021.

[40] J. Tian, K. Xu, The general position number of Cartesian products involving a factor with small diameter, Appl. Math. Comput. 403 (2021) Paper 126206.

[41] J. Tuite, E. Thomas, U. Chandran S.V., On some extremal position problems for graphs, arXiv preprint arXiv:2106.06827 (2021).

[42] Y. Yao, M. He, S. Ji, On the general position number of two classes of graphs, Open Math. 20 (2022) 1021–1029.