

# $K_{a,k}$ minors in graphs of bounded tree-width \*

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It is shown that for any positive integers  $k$  and  $w$  there exists a constant  $N = N(k, w)$  such that every  $7$ -connected graph of tree-width less than  $w$  and of order at least  $N$  contains  $K_{3,k}$  as a minor. Similar result is proved for  $K_{a,k}$  minors where  $a$  is an arbitrary fixed integer and the required connectivity depends only on  $a$ . These are the first results of this type where fixed connectivity forces arbitrarily large (nontrivial) minors.

*Key Words:* Graph theory, graph minor, connectivity.

\* This paper appeared as T. Böhme, J. Maharry, B. Mohar,  *$K_{a,k}$  minors in graphs of bounded tree-width*, J. Combin. Theory, Ser. B 86 (2002) 133–147.

<sup>†</sup> Supported in part part by the International Research Project SLO-US-007 on Graph Minors.

<sup>‡</sup> Supported in part by the Ministry of Science and Technology of Slovenia, Research Project J1–0502–0101–98, and the International Research Project SLO-US-007 on Graph Minors.

## 1. INTRODUCTION

In this paper, all graphs are finite and may have loops and multiple edges. A graph  $H$  is a *minor* of a graph  $G$ ,  $H \leq_m G$ , if  $H$  can be obtained from a subgraph of  $G$  by contracting connected subgraphs. There are many results concerning the structure of graphs that do not contain a certain graph as a minor. These excluded graphs include  $K_5$  and  $K_{3,3}$  [13],  $V_8$  [8], the 3-cube [6] and the octahedron [7]. See also [2] and [12]. There are well-known structures which guarantee a certain minor exists for large graphs. For instance, any 5-connected graph on at least 11 vertices contains the 3-cube as a minor [6]. Any 5-connected non-planar graph on at least 8 vertices contains a  $V_8$  minor [8]. In addition, there are Ramsey-type results similar to the fact that any sufficiently large connected graph contains either a  $k$ -path or a  $k$ -star. Oporowski, Oxley and Thomas [11] proved that any large 4-connected graph must have a large minor from a set of four families of graphs. Ding [3] has characterized large graphs that do not contain a  $K_{2,k}$  minor. A corollary of his result is that any large 5-connected graph contains a  $K_{2,k}$  minor.

Our results are a cross section of all of these types of results:

**THEOREM 1.1.** *For any positive integers  $k$  and  $w$  there exists a constant  $N = N(k, w)$  such that every 7-connected graph of tree-width at most  $w$  and of order at least  $N$  contains  $K_{3,k}$  as a minor.*

**THEOREM 1.2.** *There is a function  $c : \mathbb{N} \rightarrow \mathbb{N}$  such that for any  $a \geq 3$  the following holds. For any positive integers  $k$  and  $w$  there exists a constant  $N = N(k, w)$  such that every  $c(a)$ -connected graph of tree-width at most  $w$  and of order at least  $N$  contains  $K_{a,k}$  as a minor.*

Theorem 1.1 is sharp in the sense that the 7-connectivity condition cannot be relaxed. Moreover, the function  $c(a)$  in Theorem 1.2 must be at least  $2a + 1$ . These facts follow from the following construction of a family of arbitrarily large  $2a$ -connected graphs (of tree-width  $3a - 1$ ) none of which contain a  $K_{a,2a+1}$ -minor.

Let  $m$  and  $a$  be integers greater than 3. Define the graph  $N_{m,a}$  as follows. Let the vertices be indexed  $v_{x,y}$  where  $1 \leq x \leq m$  and  $1 \leq y \leq a$ . The vertex  $v_{x,y}$  is adjacent to another vertex  $v_{w,z}$  if and only if  $w \in \{x - 1, x, x + 1\}$  where  $x \pm 1$  is considered modulo  $m$ .

**PROPOSITION 1.1.** *For any integers  $a \geq 3$  and  $m \geq 3$ ,  $K_{a,2a+1} \not\leq_m N_{m,a}$ .*

*Proof.* Suppose the theorem is false for some  $a \geq 3$ . Let  $m$  be the least integer such that  $N_{m,a} \geq_m K_{a,2a+1}$ . Let the *claspers* of  $N_{m,a}$  be defined as  $CL_i = \{v_{i,y} \mid y = 1, 2, \dots, a\}$  for  $i = 1, 2, \dots, m$ .

As  $N_{m,a} \geq_m K_{a,2a+1}$ , there is a set of  $2a + 1$  connected subgraphs,  $\mathcal{S} = \{S_1, S_2, \dots, S_{2a+1}\}$ , and a set of  $a$  connected subgraphs of  $N_{m,a}$ ,  $\mathcal{T} = \{T_1, T_2, \dots, T_a\}$ , such that for every  $i, j$  there is an edge from some vertex in  $T_i$  to some vertex in  $S_j$  and such that all these subgraphs are pairwise disjoint. Assume that the  $S_i$  and  $T_i$  are chosen with  $l := \sum_{i=1}^{2a+1} |V(S_i)| + \sum_{i=1}^a |V(T_i)|$  minimum. Then it is easy to see that each of the subgraphs in  $\mathcal{S} \cup \mathcal{T}$  is a path meeting each clasp in at most one vertex. Let  $\mathcal{S}_1$  be the set of single vertex subgraphs contained in  $\mathcal{S}$ . It is easy to see that  $\mathcal{T}$  cannot contain any single vertex subgraphs.

**Claim 1:** *For every  $1 \leq i \leq m$ , there is a subgraph  $S_j \in \mathcal{S}_1$  such that  $S_j \subseteq CL_i$ .*

Suppose  $CL_i$  does not contain any of the subgraphs in  $\mathcal{S}_1$ . Then contracting a matching of size  $a$  between  $CL_i$  and  $CL_{i-1} \cup CL_{i+1}$  (indices taken modulo  $m$ ) using as many edges of  $\mathcal{S} \cup \mathcal{T}$  as possible gives a subgraph of  $N_{m-1,a}$  that still contains  $K_{a,2a+1}$  as a minor. This contradiction to the minimality of  $m$  proves the claim.

**Claim 2:** *If there is a subgraph in  $\mathcal{S}$  that contains at least two vertices, then there is a clasp that contains no member of  $\mathcal{S}_1$ .*

Suppose  $S_1$  (say) intersects  $CL_1$  and  $CL_2$ . By the minimality of  $l$ , we may assume that  $S_1 \cap CL_m = \emptyset$ . Moreover, there is a subgraph  $T_j$  that does not intersect  $CL_1 \cup CL_2 \cup CL_3$ . Otherwise, the intersection of  $S_1$  with  $CL_1$  could be removed from  $S_1$ . Therefore, a single vertex subgraph  $S_i \in \mathcal{S}_1$  contained in  $CL_2$  would not be adjacent to  $T_j$ . Hence, the clasp  $CL_2$  is as stated in the claim.

Claims 1 and 2 imply that all subgraphs in  $\mathcal{S}$  are single vertices. To complete the proof, notice that if every clasp of  $N_{m,a}$  contains one of the single vertex subgraphs of  $\mathcal{S}_1$ , then each  $T_j$  must contain at least  $m-2$  vertices in order to be adjacent to all of the subgraphs in  $\mathcal{S}$ . Hence  $|V(\mathcal{S})| + |V(\mathcal{T})| \geq |\mathcal{S}| + (m-2)|\mathcal{T}| \geq 2a+1 + (m-2)a > ma = |V(N_{m,a})|$ . This contradiction completes the proof. ■

In our proof of Theorem 1.2,  $c(3) = 7$  and  $c(a) = 264a + 1$  for  $a \geq 4$ , and we have no intention to find the best possible value for  $c(a)$ . However, the previous example shows that  $c(a)$  must be at least  $2a + 1$  for  $a \geq 3$ . It is worth remarking that our proof of Theorem 1.2 works also for  $c(a) = 3a - 1$  if we assume that the minimum degree is at least  $264a + 1$ .

## 2. BOUNDED TREE-WIDTH STRUCTURE

A *tree decomposition* of a graph  $G$  is a pair  $(T, Y)$ , where  $T$  is a tree and  $Y$  is a family  $\{Y_t \mid t \in V(T)\}$  of vertex sets  $Y_t \subseteq V(G)$ , such that the following two properties hold:

(W1)  $\bigcup_{t \in V(T)} Y_t = V(G)$ , and every edge of  $G$  has both ends in some  $Y_t$ .

(W2) If  $t, t', t'' \in V(T)$  and  $t'$  lies on the path in  $T$  between  $t$  and  $t''$ , then  $Y_t \cap Y_{t''} \subseteq Y_{t'}$ .

The *width* of a tree decomposition  $(T, Y)$  is  $\max_{t \in V(T)} (|Y_t| - 1)$ . It was shown in [11] that if a graph  $G$  has a tree decomposition of width at most  $w$  then  $G$  has a tree decomposition of width at most  $w$  that further satisfies:

(W3) For every two vertices  $t, t'$  of  $T$  and every positive integer  $k$ , either there are  $k$  disjoint paths in  $G$  between  $Y_t$  and  $Y_{t'}$ , or there is a vertex  $t''$  of  $T$  on the path between  $t$  and  $t'$  such that  $|Y_{t''}| < k$ .

(W4) If  $t, t'$  are distinct vertices of  $T$ , then  $Y_t \neq Y_{t'}$ .

(W5) If  $t_0 \in V(T)$  and  $B$  is a component of  $T - t_0$ , then  $\bigcup_{t \in V(B)} Y_t \setminus Y_{t_0} \neq \emptyset$ .

In the rest of the paper we give the proof of Theorems 1.1 and 1.2. We let  $a \geq 3$ ,  $k$ , and  $w$  be given positive integers. Let  $G$  be an  $c(a)$ -connected graph with a tree decomposition  $(T, Y)$  of width at most  $w$  that satisfies (W1)–(W5).

We will develop a structure that is similar to that used in [11]. First, we define the constants that will be used in the proofs.

$$\begin{aligned} n_5 &= r^{n_4}, \quad \text{where } r = (k-1) \binom{w+1}{a} \\ n_4 &= n_3^{w+1} \\ n_3 &= (2n_2)^p, \quad \text{where } p = 2^{w+1} \\ n_2 &= n_1^q, \quad \text{where } q = 2^{\binom{w+1}{2}} \\ n_1 &= \begin{cases} 2k(2w+3)^2 & \text{if } a = 3 \\ 2k(c(a) + 2a + 2) - 4a - 2 & \text{if } a \geq 4 \end{cases} \end{aligned}$$

We assume that  $|V(G)| = N \geq (w+1)n_5$  and that  $G$  has no  $K_{a,k}$ -minor. By (W1) we have

*Claim 2.1.*  $|V(T)| \geq n_5$ .

*Claim 2.2.* Every vertex of  $T$  has degree at most  $r = (k - 1) \binom{w+1}{a}$ .

*Proof.* Suppose  $t_0 \in V(T)$  has degree at least  $r + 1$ . Let  $\mathcal{C}$  be the set of components of  $G - Y_{t_0}$ . By (W2) and (W5), it is clear that  $|\mathcal{C}| \geq r + 1$ . For  $C \in \mathcal{C}$ , let  $X(C)$  be the set of vertices of  $Y_{t_0}$  adjacent to some vertex of  $C$ . Clearly,  $|X(C)| \geq a$  for every  $C \in \mathcal{C}$  since  $G$  is  $c(a)$ -connected and  $c(a) \geq a$ . By the Pigeonhole Principle, there is a set  $\mathcal{C}' \subseteq \mathcal{C}$  of  $k$  components for which  $\bigcap_{C \in \mathcal{C}'} X(C)$  contains  $a$  (or more) vertices of  $Y_{t_0}$ . By contracting  $B$  to a vertex for each  $B \in \mathcal{C}'$ , we see that  $G$  contains a  $K_{a,k}$  minor, a contradiction. ■

From this it follows that

*Claim 2.3.*  $T$  contains a path  $R$  of length  $|E(R)| \geq n_4$ .

The proof of the following claim can be found in [11].

*Claim 2.4.* There is a subsequence of length  $n_3$  of the vertices of  $V(R)$ ,  $r_1, r_2, \dots, r_{n_3}$ , such that for some  $s \geq 1$ ,  $|Y_{r_i}| = s$  for  $i = 1, 2, \dots, n_3$  and for every vertex of  $R$  between  $r_1$  and  $r_{n_3}$ ,  $|Y_{r_i}| \geq s$ .

From now on we replace  $R$  by the subpath from  $r_1$  to  $r_{n_3}$ . Note that because of the  $c(a)$ -connectivity and (W5),  $c(a) \leq s \leq w + 1$ .

By (W3) and Claim 2.4, there are  $s$  disjoint paths in  $G$  from  $Y_{r_1}$  to  $Y_{r_{n_3}}$ . Fix these paths, denote them by  $P_1, P_2, \dots, P_s$ , and put  $Z = P_1 \cup \dots \cup P_s$ . Since  $G$  is 3-connected, these paths can be chosen such that every  $Z$ -bridge in  $G$  is attached to at least two of the paths (cf., e.g., [4]), which we assume henceforth.

Notice that for any  $t, t' \in \{r_1, \dots, r_{n_3}\}$  and for every  $j \in \{1, \dots, s\}$  there is a unique subpath of  $P_j$  with one end in  $Y_t$  and the other end in  $Y_{t'}$ . Denote this subpath by  $P_j(t, t')$ .

The path  $P_j$  is said to be *trivial* if it consists of a single vertex, and it is said to be *everywhere nontrivial* (*almost nontrivial*) w.r.t. the sequence  $r_1, \dots, r_{n_3}$  if  $P_j(r_i, r_{i+1})$  contains at least three (respectively, at least two) vertices for each  $i = 1, \dots, n_3 - 1$ .

*Claim 2.5.* There is a subsequence  $q_1, q_2, \dots, q_{n_2}$  of  $r_1, \dots, r_{n_3}$  of length  $n_2$  such that for each  $j = 1, \dots, s$ ,  $P_j(q_1, q_{n_2})$  is either trivial or everywhere nontrivial (w.r.t. the subsequence).

*Proof.* Clearly, there is a subsequence of  $r_1, \dots, r_{n_3}$  of length  $\sqrt{n_3}$  such that the corresponding segment of  $P_1$  is either trivial or everywhere almost nontrivial with respect to the subsequence. By repeating this argument

on the subsequence for  $P_2, \dots, P_s$ , respectively, we end up with a sequence of length at least  $2n_2$  such that every path is either trivial or everywhere almost nontrivial. By taking every second element of this sequence, the required subsequence  $q_1, q_2, \dots, q_{n_2}$  is obtained.  $\blacksquare$

The paths  $P_j$  and  $P_l$  are said to be *everywhere bridge connected* (resp. *everywhere bridge disconnected*) with respect to a sequence  $p_1, \dots, p_n$  of vertices of  $R$  if for every  $i = 1, \dots, n - 1$ , there exists (resp. does not exist) a  $Z$ -bridge which has a vertex of attachment in  $P_j(p_i, p_{i+1})$  and a vertex of attachment in  $P_l(p_i, p_{i+1})$ .

*Claim 2.6.* There is a subsequence  $p_1, p_2, \dots, p_{n_1}$  of  $q_1, \dots, q_{n_2}$  of length  $n_1$  such that for every distinct pair of indices  $j, l \in \{1, \dots, s\}$ ,  $P_j(p_1, p_{n_1})$  and  $P_l(p_1, p_{n_1})$  are either everywhere bridge connected or everywhere bridge disconnected (w.r.t. the new subsequence).

*Proof.* The proof is similar to the proof of Claim 2.5 except that we have to repeat the subsequence argument  $\binom{s}{2} \leq \binom{w+1}{2}$  times.  $\blacksquare$

### 3. THE AUXILIARY GRAPH $A$

Our next goal is to examine the structure of the *auxiliary graph*  $A$  which contains information about which pairs of the paths are everywhere bridge connected. The graph  $A$  has vertex set  $V(A) = \{P_1, \dots, P_s\}$ , and the paths  $P_j$  and  $P_l$  are adjacent vertices in  $A$  if they are everywhere bridge connected w.r.t.  $p_1, \dots, p_{n_1}$  (cf. Claim 2.6).

*Claim 3.1.* Suppose that  $U \subseteq V(A)$  contains only everywhere nontrivial paths. If the subgraph of  $A$  induced by  $U$  is connected, then  $V(A) \setminus U$  contains at most  $a - 1$  vertices that are adjacent to  $U$  in  $A$ .

*Proof.* Suppose that  $P_1, \dots, P_a$  are vertices in  $V(A) \setminus U$  adjacent to  $U$  in  $A$ . Contract each path  $P_j$  ( $j = 1, \dots, a$ ) in  $G$  to a single vertex  $w_j$ . Next, for  $i = 1, 3, 5, \dots, 2k - 1$ , contract all segments  $P_j(p_i, p_{i+1})$ , where  $P_j \in U$ , and also contract all edges in bridges connecting these segments in  $G$ , to get  $k$  vertices  $z_1, z_3, \dots, z_{2k-1}$  in a minor of  $G$ . Clearly,  $n_1 \geq 2k$ , so  $z_1, z_3, \dots, z_{2k-1}$  exist. Since  $U$  is adjacent to  $P_1, \dots, P_a$  in  $A$ , it is easy to see that vertices  $w_1, \dots, w_a$  and  $z_1, z_3, \dots, z_{2k-1}$  give rise to a  $K_{a,k}$  minor of  $G$ .  $\blacksquare$

We shall apply Claim 3.1 together with the help of the following lemma.

LEMMA 3.1. *Let  $H$  be a connected graph. If  $H$  has at least  $2a^2$  vertices of degree  $\geq 3$ , then  $H$  contains a tree  $T$  with  $\geq a$  vertices of degree 1.*

*Proof.* Let  $d$  be the maximum vertex degree in  $H$ , and let  $v_0$  be a vertex of degree  $d$ . If  $d \geq a$ , then  $T$  is the star centered at  $v_0$ . So, suppose that  $d < a$ . Then it is sufficient to prove the following. Assuming that  $H$  has at least  $2a^2 - (d-1)^2$  vertices of degree  $\geq 3$ , we shall prove by induction on  $a-d$  that the tree  $T$  exists. Let  $N_1$  be the set of all vertices of degree  $\geq 3$  which can be reached from  $v_0$  on paths whose internal vertices all have degree 2. Then  $1 \leq |N_1| \leq d$ . Let  $N_2$  be the “second neighborhood” of  $v_0$ , consisting of vertices of degree  $\geq 3$  which are not in  $N_1 \cup \{v_0\}$  and which can be reached from  $v_0$  on paths for which exactly one internal vertex has degree  $\geq 3$ . Similarly, let  $N_3$  be the “third neighborhood” of  $v_0$ . Then  $1 \leq |N_2| \leq d(d-1)$  and  $|N_3| \geq 1$  since  $H$  is connected and  $2a^2 - (d-1)^2 > 1 + d + d(d-1) \geq 1 + |N_1| + |N_2|$ . Let  $v_3 \in N_3$ , and let  $W$  be a path from  $v_0$  to  $v_3$  which contains precisely two other vertices of degree  $\geq 3$ . Now, contract  $W$  to a vertex  $\tilde{v}_0$  and remove possible parallel edges. Denote the resulting graph by  $\tilde{H}$ . If a vertex of  $\tilde{H}$  has degree smaller than in  $H$ , then it was adjacent to two (or three) vertices of  $W$ . This implies that  $\tilde{H}$  has at least  $2a^2 - (d-1)^2 - (2d-1) = 2a^2 - ((d+1)-1)^2$  vertices of degree  $\geq 3$ . Since  $v_0$  and  $v_3$  have no common neighbors,  $\tilde{v}_0$  is its vertex of maximum degree  $\geq d+1$ . By the induction hypothesis,  $\tilde{H}$  contains a tree  $\tilde{T}$  with at least  $a$  vertices of degree 1. Clearly,  $\tilde{T}$  gives rise to the required tree  $T$  in  $H$ . ■

At least one of the paths is everywhere nontrivial, say  $P_1$ . Let  $A_1$  be the induced subgraph of  $A$  on the everywhere nontrivial paths. Let  $A_0$  be the induced subgraph of  $A$  consisting of the connected component of  $A_1$  containing  $P_1$  together with (at most  $a-1$ ) trivial paths adjacent to that component.

From now on we shall assume that  $G$  is  $c(a)$ -connected, where  $c(3) = 7$  and  $c(a) = 264a + 1$  for  $a \geq 4$ .

*Claim 3.2.*  $A_0 \cap A_1$  has at least  $\lceil \frac{c(a)-a+1}{2} \rceil$  vertices. If  $a = 3$ ,  $A_0$  is isomorphic to a path or a cycle on at least four vertices. If  $a \geq 4$ , then every vertex of  $A_0 \cap A_1$  has degree at most  $a-1$  and at most  $2a^2$  of these vertices have degree more than 2 in  $A_0 \cap A_1$ .

*Proof.* Let  $U = V(A_0 \cap A_1)$ ,  $x = |U|$ , and  $y = |V(A_0)| - x$ . By Claim 3.1 we see that  $y \leq a-1$ . Since the  $2x+y$  endvertices of the paths in  $A_0$  in  $Y_{p_1}$  and  $Y_{p_3}$  separate the graph  $G$ , we have  $2x+y \geq c(a)$ . This implies that  $x \geq (c(a) - a + 1)/2$ , and proves the first part of the claim.

By Claim 3.1, every vertex in  $A_0 \cap A_1$  has degree at most  $a - 1$  in  $A$ . If  $a = 3$ , this implies that  $A_0 \cap A_1$  is a path or a cycle, and the trivial paths in  $V(A_0)$  can be adjacent only to vertices of degree  $\leq 1$  in  $A_0 \cap A_1$ . This and Claim 3.1 imply that  $A_0$  is a path or a cycle. If  $|V(A_0)| \leq 3$ , then the endpoints of the paths in  $V(A_0)$  would give a  $\leq 6$ -separator in  $G$ .

Suppose now that  $a \geq 4$ . By Claim 3.1 every vertex of  $A_0 \cap A_1$  has degree at most  $a - 1$ . Suppose that there are more than  $2a^2$  vertices of degree  $\geq 3$ . By Lemma 3.1,  $A_0 \cap A_1$  contains a tree  $T$  with  $\geq a$  vertices of degree 1. Let  $U$  be the set of vertices of degree  $\geq 2$  in  $T$ . The subgraph of  $A$  induced by  $U$  is connected, and Claim 3.1 yields a contradiction. This completes the proof.  $\blacksquare$

Denote by  $Z'(i)$  the union of  $P_j(p_i, p_{i+1})$  where  $P_j \in V(A_0)$ ,  $i = 1, 2, \dots, n_1 - 1$ . Let  $Z_i$  be the subgraph of  $G$  obtained by taking the union of  $Z'(i)$  and all those  $Z$ -bridges  $B$  that have all vertices of attachment in  $Z'(i)$  such that there is no  $i' < i$  for which  $B$  would have all its vertices of attachment in  $Z'(i')$ .

#### 4. FINDING $K_{3,K}$ MINORS

In this section we consider the case when  $a = 3$  since the best possible connectivity 7 requires more elaborate techniques than the general case treated in the next section. For  $i = 1, 2, \dots, n_1 - 2w - 2$ , let  $H_i = \bigcup_{k=0}^{2w} Z_{i+k}$ . Let  $R, R' \in V(A_0)$  be paths which are adjacent in  $A_0$ . For  $i = 1, 2, \dots, n_1 - 2w - 2$  define the graph  $D_i = D_i(R, R')$  as follows. First, take  $S = (R \cup R') \cap H_i$  together with all  $Z$ -bridges in  $H_i$  that have vertices of attachment on  $R$  and on  $R'$ . Finally, add two edges  $e_1, e_2$ , where  $e_1$  joins the “left” endvertices,  $\lambda$  in  $R \cap H_i$  and  $\lambda'$  in  $R' \cap H_i$ , and  $e_2$  joins the “right” endvertices,  $\rho$  and  $\rho'$ , of these two paths. Then  $S + e_1 + e_2 =: C$  is a cycle in  $D_i$ . If  $R$  ( $R'$ ) is everywhere trivial, then  $\lambda = \rho$  ( $\lambda' = \rho'$ ).

*Claim 4.1.* Suppose that  $a = 3$ . Then for every  $i$ , there are adjacent vertices  $R, R'$  of  $A_0$  such that  $D_i(R, R')$  has no embedding in the plane where the vertices  $\lambda, \lambda', \rho', \rho$  would lie on the outer face in the prescribed order.

*Proof.* Suppose that  $H_i$  is a planar graph. Let  $v_j$  be the number of vertices of degree  $j$  in  $H_i$ . By Euler’s formula and standard counting arguments it follows that

$$L := \sum_{j \geq 0} (6 - j)v_j \geq 12. \quad (1)$$



Observe that  $H_i$  has at most  $2s$  vertices of degree  $\leq 6$  since the minimum degree in  $G$  is at least 7 (by the 7-connectivity of  $G$ ). On the other hand, since at least three of the paths in  $H_i$  are nontrivial, these paths contain at least  $3(2(2w+1)-1) = 12w+3$  vertices of degree  $\geq 7$  in  $H_i$ . Therefore,

$$L \leq 6 \cdot 2s - (12w+3) \leq 12(w+1) - 12w - 3 = 9.$$

This contradiction to (1) shows that  $H_i$  is not planar. Recall that  $A_0$  is a path or a cycle on at least 4 vertices,  $R_1, \dots, R_d$ ,  $d \geq 4$ . This implies, in particular, that no  $Z$ -bridge in  $H_i$  is attached to more than two of the paths (otherwise, there would be a 3-cycle in  $A_0$ , and so  $A_0$  would be equal to the 3-cycle). Moreover, if every  $D_i(R_j, R_{j+1})$  ( $j = 1, \dots, d$ , indices taken modulo  $d$ ) has an embedding in the plane with the corresponding cycle  $C_j$  being the outer cycle, then  $\bigcup_{j=1}^d D_i(R_j, R_{j+1}) \supseteq H_i$  would be planar as well, contrary to the above. Hence, there is an index  $j$  such that  $D_i(R_j, R_{j+1})$  has no such embedding. Since there are no local  $Z$ -bridges,  $D_i(R_j, R_{j+1})$  neither has an embedding in the plane where the vertices  $\lambda, \lambda', \rho', \rho$  are on the outer face in the prescribed order. ■

We shall need a result about crossing paths from from [9]. A *separation* of a graph  $G$  is a pair  $(A, B)$  of subgraphs with  $A \cup B = G$  and  $E(A \cap B) = \emptyset$ , and its *order* is  $|V(A \cap B)|$ . By a *society* we mean a pair  $(G, \Omega)$ , where  $G$  is a graph and  $\Omega$  a cyclic permutation of a subset  $\overline{\Omega}$  of  $V(G)$ . A *cross* in  $(G, \Omega)$  is a pair of disjoint paths in  $G$  with ends  $s_1, t_1$  and  $s_2, t_2$ , respectively, all in  $\overline{\Omega}$ , such that  $s_1, s_2, t_1, t_2$  occur in  $\Omega$  in that order (but not necessarily consecutive). The following formulation of a theorem of Robertson and Seymour [9] appears in [10].

**THEOREM 4.1** (Robertson and Seymour). *Let  $(G, \Omega)$  be a society such that there is no separation  $(A, B)$  of  $G$  of order  $\leq 3$  with  $\overline{\Omega} \subseteq V(A) \neq V(G)$ . Then the following are equivalent:*

- (a) *There is no cross in  $(G, \Omega)$ .*
- (b)  *$G$  can be drawn in a disc with the vertices in  $\overline{\Omega}$  drawn on the boundary of the disc in order given by  $\Omega$ .*

*Claim 4.2.* If  $D_i(R, R')$  is nonplanar, then one of the following holds:

- (a)  $D_i(R, R')$  contains disjoint paths  $Q_1, Q_2$  connecting  $\lambda$  with  $\rho'$  and  $\lambda'$  with  $\rho$ , respectively.
- (b)  $D_i(R, R')$  contains a path  $Q$  (resp.,  $Q'$ ) disjoint from  $R'$  (resp.,  $R$ ) which connects  $\lambda$  and  $\rho$  (resp.,  $\lambda'$  and  $\rho'$ ) such that after replacing  $R$  (resp.,  $R'$ ) by  $Q$  (resp.,  $Q'$ ), there is a  $Z$ -bridge in  $H_i$  which is attached to more than two of the paths  $P_1, \dots, P_s$ .

*Proof.* Let  $H = D_i(R, R')$ . Let  $C$  be the cycle of  $H$  defined before Claim 4.1. Let  $\bar{\Omega}$  be the set of vertices of  $C$  which are incident with an edge in  $E(G) \setminus E(H)$ . The cyclic order of  $\bar{\Omega}$  on  $C$  defines the society  $(H, \Omega)$ . Since  $G$  is 4-connected and no vertex in  $V(H) \setminus \bar{\Omega}$  is incident with an edge in  $E(G) \setminus E(H)$ , there is no separation  $(A, B)$  of  $H$  of order  $\leq 3$  with  $\bar{\Omega} \subseteq V(A) \neq V(G)$ . Since  $H$  is nonplanar, Theorem 4.1 implies that there is a cross  $R_1, R_2$  in  $(H, \Omega)$ . Let  $\alpha_i, \beta_i$  be the endvertices of  $R_i$  ( $i = 1, 2$ ). We may assume that:

- (i) None of the vertices  $\lambda, \lambda', \rho, \rho'$  is an internal vertex of  $R_1$  or  $R_2$ .

Subject to (i) choose the cross  $R_1, R_2$  such that

- (ii)  $\{\alpha_1, \alpha_2, \beta_1, \beta_2\}$  contains as many vertices in  $\{\lambda, \lambda', \rho, \rho'\}$  as possible and, subject to (i) and (ii)

- (iii) as few edges in  $E(H) \setminus E(R \cup R')$  as possible.

If  $\lambda, \lambda', \rho, \rho'$  are all endvertices of  $R_1, R_2$ , then we have (a). Hence we may assume that  $\lambda$  is not an endvertex of  $R_1, R_2$ . If  $R \cap (R_1 \cup R_2) \neq \emptyset$ , let  $v$  be the first vertex of  $R_1 \cup R_2$  on  $R$  (starting at  $\lambda$  towards  $\rho$ ). We may assume that  $v \in V(R_1)$ . Let  $R_1 = R'_1 \cup R''_1$  where  $V(R'_1) \cap V(R''_1) = \{v\}$ . By replacing one of the segments  $R'_1$  or  $R''_1$  in  $R_1$  by a segment from  $v$  to  $\lambda$  on  $R$ , a new cross is obtained which contradicts (ii) or (iii), except when  $R'_1$  or  $R''_1$  is the segment of  $R$  from  $v$  to  $\rho$ . In particular, three of the endvertices of  $R_1, R_2$  are on  $R'$ . The above proof implies that  $\lambda'$  and  $\rho'$  are the endvertices of the paths. Since  $R_1, R_2$  cross,  $R_1$  joins a vertex  $x \in V(R') \setminus \{\lambda', \rho'\}$  with  $\rho$ , and  $R_2$  joins  $\lambda'$  and  $\rho'$ , where  $R_2$  is disjoint from  $R$ . It is easy to see, that this gives (b).

Suppose now that  $R \cap (R_1 \cup R_2) = \emptyset$ . Condition (ii) implies that  $\lambda'$  and  $\rho'$  are the endvertices of  $R_1$  and  $R_2$ , respectively. There is a  $C$ -bridge  $B$  in  $H$  such that  $E(R_1 \cup R_2) \cap E(B) \neq \emptyset$ . Since  $B$  is not a local bridge, it is attached to  $R$  as well. Therefore, there is a path  $L$  in  $B$  from  $R$  to  $R_1 \cup R_2$  (say to  $R_2$ ) which is internally disjoint from  $C \cup R_1 \cup R_2$ . Let  $y$  be the vertex of  $R_1$  which is as close as possible to  $\rho'$  on  $R'$ . Let  $R'_2$  be the segment of  $R_2$  from  $R_2 \cap L$  to the end of  $R_2$  distinct from  $\rho'$ . By (iii),  $R'_2$  is disjoint from the segment  $Q''$  of  $R'$  from  $y$  to  $\rho'$ . Therefore, the path  $Q'$  composed of the segment of  $R_1$  from  $\lambda'$  to  $y$  and  $Q''$  can be taken as the path  $Q'$  in (b). Note that, after replacing  $R'$  by  $Q'$ , the  $Z$ -bridge containing  $L \cup R'_2$  will be attached to at least three paths in  $\{P_1, \dots, P_s\}$ . ■

We are ready to complete the proof of Theorem 1.1. Suppose that  $a = 3$  and that  $A_0$  is a path or a cycle on consecutive vertices  $R_1, \dots, R_d$ , where  $4 \leq d \leq w + 1$ . Let  $D_i^j = D_i(R_j, R_{j+1})$ ,  $j = 1, \dots, d$ . We shall only

consider the indices  $i$  of the form  $i = 1 + t(2w + 2)$ ,  $t = 0, 1, \dots$ , and we call them *admissible indices*.

Let us first assume that the case (b) of Claim 4.2 occurs less than  $2kd$  times at admissible indices  $i$ . Since there are at least  $4kd$  admissible indices, Claim 4.2(a) implies that there is an index  $j \in \{1, \dots, d\}$ , and there are admissible indices  $1 \leq i_1 < i_2 < \dots < i_k \leq n_1 - 2w - 2$  such that

- (i) each of  $D_{i_1}^j, D_{i_2}^j, \dots, D_{i_k}^j$  contains paths as stated in Claim 4.2(a), and
- (ii) for  $l = 1, \dots, k - 1$ ,  $i_{l+1} - i_l \geq 2w + 2$ .

We can exchange the segments of the paths  $R_j$  and  $R_{j+1}$  in  $H_{i_l}$  by the two paths  $Q_1, Q_2$  of Claim 4.2(a). In this way the new paths in  $H_{i_l} \cup Z_{i_l+2w+2}$  would no longer satisfy the condition of Claim 3.1. Namely, if  $R_j$  and  $R_{j+1}$  have degrees  $d_1, d_2$  in  $A_0$ , then they would be everywhere bridge connected (w.r.t. the sequence  $p_{i_1-1}, p_{i_2-1}, \dots, p_{i_k-1}$ ) with  $d_1 + d_2 - 1$  other paths. If  $d_1 = d_2 = 2$ , this gives a  $K_{3,k}$  minor in the same way as in the proof of Claim 3.1 (since one of  $R_j$  or  $R_{j+1}$  is everywhere nontrivial). If  $d_1 = 1$  (say), then the path  $R_{j+2}$  has degree 2 in  $A_0$  by Claim 3.2 and (in addition to  $R_{j+3}$ ) it becomes everywhere bridge connected to the two new paths (w.r.t. the sequence  $p_{i_1-1}, p_{i_2-1}, \dots, p_{i_k-1}$ ). It is easy to see from the definition of  $A_0$  that  $R_{j+2}$  cannot be trivial, so the proof of Claim 3.1 applies again.

Let us now assume that the case (b) of Claim 4.2 occurs  $2kd$  or more times (for admissible indices  $i$ ). Then there is an index  $j \in \{1, \dots, d\}$ , and there are admissible indices  $1 \leq i_1 < i_2 < \dots < i_k \leq n_1 - 2w - 2$  such that

- (i) each of  $D_{i_1}^j, D_{i_2}^j, \dots, D_{i_k}^j$  contains a path  $Q$  (or each of  $D_{i_1}^j, D_{i_2}^j, \dots, D_{i_k}^j$  contains a path  $Q'$ ) as stated in Claim 4.2(b), and
- (ii) for  $l = 1, \dots, k - 1$ ,  $i_{l+1} - i_l \geq 2w + 2$ .

For any  $D_{i_l}^j$  we replace the segment of  $R_j$  (resp.,  $R_{j+1}$ ) by the corresponding path  $Q$  (resp.,  $Q'$ ) such that there is a  $Z$ -bridge (where  $Z$  is defined as the union of the new paths) attached to  $R_j, R_{j+1}$ , and  $R_{j+2}$  (or  $R_{j-1}$ ). We may assume that  $k$  of these bridges,  $B_1, \dots, B_k$  are attached to  $R_j, R_{j+1}$ , and  $R_{j+2}$ . Now, there is a  $K_{3,k}$ -minor obtained by contracting  $R_j, R_{j+1}, R_{j+2}$  into single vertices and adding paths in  $B_1, \dots, B_k$  to these vertices. This completes the proof of Theorem 1.1.

## 5. FINDING $K_{A,K}$ MINORS FOR $A \geq 4$

Suppose now that  $a \geq 4$  and  $c(a) = 264a + 1$ . Let  $r = 2c(a) + 2$ . For  $i = 1, 2, \dots, n_1 - r$ , let  $H_i = \bigcup_{j=0}^{r-1} Z_{i+j}$ . We also write  $S_i = Y_{p_i}$ .

*Claim 5.1.* For every  $1 \leq i \leq n_1 - r$ , the average degree of vertices in  $H_i$  is at least  $c(a) - \frac{1}{2}$ .

*Proof.* Every vertex of  $G$  has degree at least  $c(a)$ . Let  $s_0 = |V(A_0 \cap A_1)|$  be the number of everywhere nontrivial paths in  $V(A_0)$ . Then

$$|V(H_i)| \geq s_0(2r + 1) > 4s_0c(a). \quad (2)$$

Each trivial path in  $V(A_0)$  is everywhere bridge connected to some nontrivial path. Hence, the degree of the corresponding vertex in  $H_i$  is at least  $r/2 \geq c(a)$ . Only the ends of nontrivial paths can have degree less than  $c(a)$  in  $H_i$ . This fact and inequality (2) imply that

$$2|E(H_i)| \geq c(a)(|V(H_i)| - 2s_0) \geq (c(a) - \frac{1}{2})|V(H_i)|.$$

This completes the proof.  $\blacksquare$

A graph  $L$  is said to be  $q$ -linked if it has at least  $2q$  vertices and for any ordered  $q$ -tuples  $(s_1, \dots, s_q)$  and  $(t_1, \dots, t_q)$  of  $2q$  distinct vertices of  $L$ , there exist pairwise disjoint paths  $P_1, \dots, P_q$  such that for  $i = 1, \dots, q$ , the path  $P_i$  connects  $s_i$  and  $t_i$ . Such collection of paths is called a *linkage* of  $(s_1, \dots, s_q)$  and  $(t_1, \dots, t_q)$ .

*Claim 5.2.* For every  $1 \leq i \leq n_1 - r$ , there exists a subgraph  $L_i$  of  $H_i$  which is  $3a$ -linked.

*Proof.* Mader [5] proved that every graph of average degree at least  $4c$  contains a  $c$ -connected subgraph. Therefore, since  $H_i$  has average degree at least  $c(a) - 1 \geq 264a$ ,  $H_i$  contains a  $66a$ -connected subgraph  $L_i$ . Bollobás and Thomason [1] have shown that every  $22t$ -connected graph is  $t$ -linked. Hence, the graph  $L_i$  is  $3a$ -linked.  $\blacksquare$

We will now construct  $a$  disjoint paths  $\mathcal{P}_1^\circ, \dots, \mathcal{P}_a^\circ$  by routing the paths  $P_1, \dots, P_s$  through  $L_i$  in at least  $k$  pairwise disjoint subgraphs  $H_i$ . In each graph  $L_i$ , there will also be an extra vertex linked to each of the  $a$  paths. Contracting these paths will then give a  $K_{a,k}$ -minor in  $G$ .

*Claim 5.3.* In  $H_i$ , there exist  $2a$  pairwise disjoint paths,  $Q_1^{(i)}, \dots, Q_a^{(i)}$  and  $Q'_1{}^{(i)}, \dots, Q'_a{}^{(i)}$  such that the following hold:

- (a) For  $l = 1, 2, \dots, a$ , the path  $Q_l^{(i)}$  starts in  $L_i$  and ends in  $S_{i+r}$ .
- (b) For  $l = 1, 2, \dots, a$ , the path  $Q'_l{}^{(i)}$  starts in  $S_i$  and ends in  $L_i$ .
- (c) Every path  $Q_l^{(i)}$  and  $Q'_l{}^{(i)}$  ( $l = 1, 2, \dots, a$ ) has only its endvertices in  $S_i \cup S_{i+r} \cup V(L_i)$ .

*Proof.* Let  $\Pi_0 = V(A_0) \setminus V(A_1)$  be the set of vertices of  $H_i$  corresponding to the trivial paths in  $A_0$ . Let  $\mathcal{W} = \{W_1, \dots, W_{2a}\}$  be a set of  $2a$  pairwise disjoint paths joining  $V(L_i)$  with  $S_i \cup S_{i+r}$  such that:

- (1)  $W_l \subseteq H_i - \Pi_0$  for every  $l = 1, 2, \dots, 2a$ .
- (2) The number of edges in  $\bigcup_{l=1}^{2a} E(W_l) \setminus \bigcup_{j=0}^{r-1} E(Z'(i+j))$  is minimum.
- (3) Subject to (2), if  $n_L$  is the number of paths  $W_l$  ending in  $S_i$ , and  $n_R$  is the number of paths  $W_l$  ending in  $S_{i+r}$ ,  $|n_L - n_R|$  is minimum.

Disjoint paths satisfying (1) exist by large connectivity: Since  $c(a) \geq 3a - 1$ , and  $|V(L_i)| > 3a$ , and  $|S_i \cup S_{i+r}| \geq 3a - 1$ , there exist  $3a - 1$  disjoint paths from  $V(L_i)$  to  $S_i \cup S_{i+r+1}$  by Menger's theorem. Since there are at most  $a - 1$  vertices in  $\Pi_0$ , the removal of those paths which intersect  $\Pi_0$  leaves at least  $2a$  paths satisfying condition (1).

If at least two paths of  $\mathcal{W}$  intersect a path  $P_j$ , then let  $W$  and  $W'$  be the paths that intersect  $P_j$  as close as possible (on  $P_j$ ) to  $S_i$  and  $S_{i+r}$ , respectively. If  $W = W'$ , suppose that the intersection  $u$  of  $W$  with  $P_j$  nearest  $S_i$  (say) comes before the intersection nearest  $S_{i+r}$ . By (2),  $W$  ends at  $S_i$ , i.e., its segment from  $u$  to its end coincides with the segment  $P_j(u, S_i)$  of  $P_j$ . This shows that  $W \neq W'$ . Then the path  $W$  (resp.  $W'$ ) must end at  $S_i$  (resp.  $S_{i+r}$ ) by (2).

Suppose that precisely one path, say  $W \in \mathcal{W}$ , intersects a path  $P_j$ . In this case we can elect to have  $W$  ending at  $P_j \cap S_i$  or at  $P_j \cap S_{i+r}$  by following the path  $P_j$ . This implies that the value  $|n_L - n_R|$  in (3) can be made to be zero. Then  $n_L = n_R = a$ .

Now let the  $a$  paths in  $\mathcal{W}$  that end in  $S_i$  be called  $Q_1^{(i)}, Q_2^{(i)}, \dots, Q_a^{(i)}$  and the  $a$  paths in  $\mathcal{W}$  that end in  $S_{i+r}$  be called  $Q_1^{(i)}, Q_2^{(i)}, \dots, Q_a^{(i)}$ . It is easy to see that (c) may be requested. This completes the proof.  $\blacksquare$

Let  $T$  be a spanning tree of  $A_0 \cap A_1$ . By Claim 3.2,  $|V(T)| \geq a$ . This implies the following claim.

*Claim 5.4.* There are vertices  $t_1, t_2, \dots, t_a$  of  $T$  such that for  $l = 1, 2, \dots, a$ , the vertex  $t_l$  is a leaf of the subtree  $T \setminus \{t_1, \dots, t_{l-1}\}$ .

For each  $i = 1, 2, \dots, n_1 - r$  and each  $l = 1, 2, \dots, a$ , let  $J_l^{(i)} \in \{P_1, \dots, P_s\}$  be the vertex of  $T$  such that  $Q_l^{(i)}$  ends up on the corresponding path in  $G$ . Choose an enumeration of  $Q_1^{(i)}, Q_2^{(i)}, \dots, Q_a^{(i)}$  such that, for  $l = 1, 2, \dots, a$ , the distance from  $J_l^{(i)}$  to  $t_l$  in  $T$  is minimum (where smaller values of  $l$  have preference over the larger values).

Choose a similar enumeration of  $Q_1^{(i)}, \dots, Q_a^{(i)}$ .

Define  $\alpha = r + 4a + 2$  and for  $t = 1, \dots, k$  set  $i_t = 1 + (t - 1)\alpha$ . Observe that  $i_k = n_1 - r$ .

To construct the path  $\mathcal{P}_l^\circ$ , we first link  $Q_l^{(i_t)}$  to  $Q_l^{(i_{t+1})}$  for every  $t = 1, \dots, k-1$ . Then each  $Q_l^{(i_t)}$  is linked to  $Q_l^{(i_t)}$  inside  $L_{i_t}$  ( $t = 1, \dots, k$ ). We do this as described below.

Let  $i' = i + \alpha$ . Link  $Q_l^{(i)}$  with  $Q_l^{(i')}$  as follows: Follow the path  $J_l^{(i)}$  from  $J_l^{(i)} \cap S_{i+r}$  through  $2l$  segments to the separator  $S_{i+r+2l}$ . Continue the path within  $Z_{i+r+2l}$  to the path  $t_l$ . This can be done by following the bridges between paths corresponding to the path in the spanning tree  $T$  from  $J_l^{(i)}$  to  $t_l$ .

Construct a similar path from  $Q_l^{(i')}$  to  $t_l$  using bridges in  $Z_{i'-2l}$ . Then connect these paths along  $t_l$ , and denote by  $P_l^i$  the resulting path joining  $Q_l^{(i)}$  with  $Q_l^{(i')}$ .

*Claim 5.5.* The constructed paths  $P_l^i$  ( $l = 1, \dots, a$ ) are pairwise disjoint.

*Proof.* Consider two of the paths, say  $P_l^i$  and  $P_m^i$ , where  $l < m$ . There are four possibilities where these two paths may intersect:

- (1)  $P_l^i$  intersects  $J_m^{(i)}$  inside  $Z_{i+r+2l}$ : This is not possible since  $J_m^{(i)}$  would then be closer to  $t_l$  in  $T$ , and the path  $Q_m^{(i)}$  would be indexed before  $Q_l^{(i)}$ .
- (2)  $P_m^i$  intersects  $t_l$  inside  $Z_{i+r+2m}$ : This is not possible since  $t_l$  is a leaf in  $T \setminus \{t_1, \dots, t_{l-1}\}$ .

The remaining cases, when  $P_l^i$  intersects  $P_m^i$  inside  $Z_{i'-2l}$  or inside  $Z_{i'-2m}$  (respectively) are handled similarly. This completes the proof.  $\blacksquare$

Let  $v_l$  be the vertex of  $Q_l^{(i)}$  in  $L_i$ , and let  $u_l$  be the vertex of  $Q_l^{(i)}$  in  $L_i$ . Choose  $u'_l$  to be a neighbor of  $u_l$  in  $L_i \setminus \{v_1, \dots, v_a, u_1, \dots, u_a\}$ . Since  $L_i$  is  $3a$ -linked, the minimum degree of  $L_i$  is at least  $3a$ , so such neighbors exist. The vertices  $u'_l$  may even be chosen so that they are pairwise distinct. Let  $v'_1 = u'_1$ , and let  $v'_2, \dots, v'_a$  be distinct neighbors of  $v'_1$  in  $L_i$ . We may assume that if  $v'_\alpha = u'_\beta$ , then  $\alpha = \beta$ .

Since  $L_i$  is  $2a$ -linked, there is a linkage from  $(v_1, \dots, v_a, v'_1, \dots, v'_a)$  to  $(u_1, \dots, u_a, u'_1, \dots, u'_a)$ . The resulting paths joining  $v_l$  and  $u_l$  ( $l = 1, \dots, a$ ) are used to link  $Q_l^{(i)}$  and  $Q_l^{(i)}$  inside  $L_i$ , for  $i \in \{i_1, \dots, i_k\}$ . Together with the paths  $P_l^i$ ,  $i \in \{i_1, \dots, i_{k-1}\}$ , this determines the path  $\mathcal{P}_l^\circ$ . On the other hand, the paths in the linkage from  $(v'_1, \dots, v'_a)$  to  $(u'_1, \dots, u'_a)$  are disjoint from  $\mathcal{P}_1^\circ, \dots, \mathcal{P}_a^\circ$  and can be used to link  $v'_1$  to each of these paths.

Now, it can be shown that  $G$  contains a  $K_{a,k}$  minor: For each  $l = 1, \dots, a$ , contract the path  $\mathcal{P}_l^\circ$  to a single vertex. For  $i \in \{i_1, \dots, i_k\}$ , the vertex  $v'_1 \in V(L_i)$  is joined to  $u'_1, \dots, u'_a$  and hence to each of the  $a$  paths  $\mathcal{P}_1^\circ, \dots, \mathcal{P}_a^\circ$ . Since this is repeated  $k$  times, we get a  $K_{a,k}$  minor in  $G$ .

The proof of Theorem 1.2 is complete.

## 6. CONCLUSION

Our more recent results show that the condition on bounded tree-width in Theorem 1.1 can be removed. The authors plan a second paper in which the large tree-width case is handled. This will prove the following, which was conjectured independently by Ding [3] and the authors:

*There is a function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that any 7-connected graph on at least  $f(k)$  vertices contains a  $K_{3,k}$  minor.*

It seems reasonable to the authors that this result can be extended to  $K_{4,k}$ -minors and possibly even to  $K_{a,k}$ -minors. The logical conjectures would be the following:

*Conjecture 6.1.* There is a function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that any 9-connected graph on at least  $f(k)$  vertices contains a  $K_{4,k}$  minor.

*Conjecture 6.2.* There are functions  $f : \mathbb{N} \rightarrow \mathbb{N}$  and  $c : \mathbb{N} \rightarrow \mathbb{N}$  such that any  $c(a)$ -connected graph on at least  $f(k)$  vertices contains a  $K_{a,k}$  minor.

Our final remark is that the sequence of graphs  $K_{a,k}$ , where  $a$  is fixed and  $k$  tends to infinity, is essentially the only family of graphs for which a result like our Theorem 1.2 holds. More precisely:

**THEOREM 6.1.** *Let  $c$  and  $w \geq c$  be positive integers, and let  $H_k$  ( $k \geq 1$ ) be a sequence of graphs such that  $\lim_{k \rightarrow \infty} |V(H_k)| = \infty$ . Suppose that for any positive integer  $k$  there exists an integer  $N(k)$  such that every  $c$ -connected graph of tree-width  $\leq w$  and of order at least  $N(k)$  contains  $H_k$  as a minor. Then  $H_k \leq_m K_{c,N(k)}$  for  $k \geq 1$ .*

*Proof.* Clearly, the graph  $K_{c,N(k)}$  is  $c$ -connected and has tree-width  $c \leq w$ . By the assumption on the family  $H_k$ ,  $K_{c,N(k)}$  contains  $H_k$  as a minor. ■

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